

C语言拾遗：机制和实践

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讲前提醒

Token已发送，PA0截止时间：

- * 2025年9月21日23:59:59 (以此 deadline 计按时提交bonus)

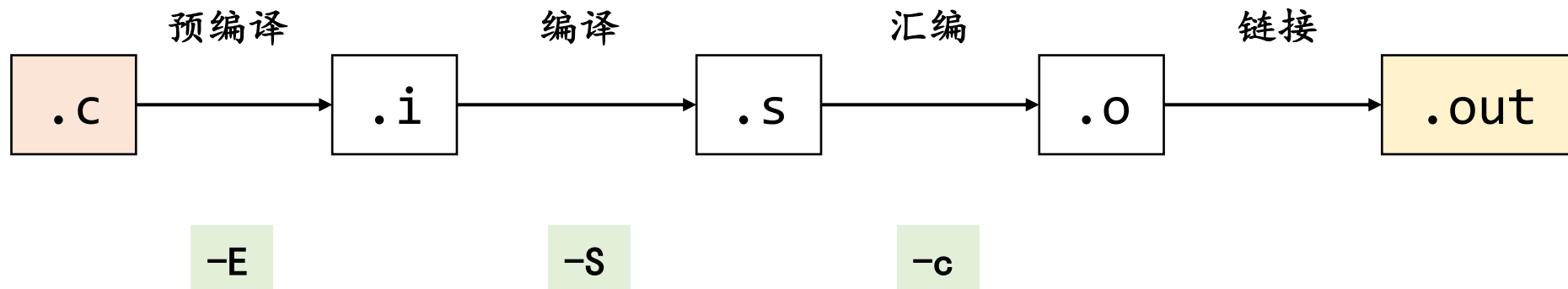
PA1已悄悄发布：（OJ已开启PA1评测）

- * PA 1.1: 2025.10.5 (此为建议的不计分 deadline)

- * PA 1.2: 2025.10.12 (此为建议的不计分 deadline)

- * PA 1.3: 2025.10.19 23:59:59 (以此 deadline 计按时100%提交)

从源代码到可执行文件



预编译

(先行剧透本学期的主要内容)

.c → 预编译 → .i → 编译 → .s → 汇编 → .o → 链接 → .out

X-Macros

- 宏展开：通过复制/粘贴改变代码的形态
 - 反复粘贴，直到没有宏可以展开为止
- 例子：X-macro

```
$ ./a.out  
Hello, Tom!  
Hello, Jerry!  
Hello, Tyke!  
Hello, Spike!
```

```
#define NAMES(X) \  
    X(Tom) X(Jerry) X(Tyke) X(Spike)  
  
int main() {  
    #define PRINT(x) puts("Hello, " #x "!!");  
    NAMES(PRINT)  
}
```

PRINT(TOM) PRINT(Jerry) PRINT(Tyke) PRINT(Spike)

X-Macros

```
#define NAMES(X) \  
    X(Tom) X(Jerry) X(Tyke) X(Spike)
```

```
int main() {  
    #define PRINT(x) puts("Hello, " #x "!");  
    NAMES(PRINT)  
    #define PRINT2(x) puts("Goodbye, " #x "!");  
    NAMES(PRINT2)  
}
```

```
PRINT(TOM) PRINT(Jerry) PRINT(Tyke) PRINT(Spike)
```

```
PRINT2(TOM) PRINT2(Jerry) PRINT2(Tyke) PRINT2(Spike)
```

```
$ ./a.out  
Hello, Tom!  
Hello, Jerry!  
Hello, Tyke!  
Hello, Spike!  
Goodbye, Tom!  
Goodbye, Jerry!  
Goodbye, Tyke!  
Goodbye, Spike!
```

PA中有趣的预编译

- PA头文件遵循标准头文件结构，避免被重复引用

- in nemu/include/isa.h

```
#ifndef __ISA_H__  
#define __ISA_H__  
.....  
#endif
```

- in nemu/include/common.h
 - 实际上是include哪一些头文件？
 - CONFIG_TARGET_AM是否有定义，在哪里？
 - make menuconfig做了什么？

```
#ifdef CONFIG_TARGET_AM  
#include <klib.h>  
#else  
#include <assert.h>  
#include <stdlib.h>  
#endif
```

PA中有趣的预编译

```
#ifdef cplusplus
extern "C" {
#endif

// ----- TRM: Turing Machine -----
extern Area heap;
void putch (char ch);
void halt (int code) __attribute__((__noreturn__));

// ----- IOE: Input/Output Devices -----
bool ioe_init (void);
void ioe_read (int reg, void *buf);
void ioe_write (int reg, void *buf);
#include "amdev.h"

// ----- CTE: Interrupt Handling and Context Switching -----
bool cte_init (Context *(*handler)(Event ev, Context *ctx));
void yield (void);
bool ienabled (void);
void iset (bool enable);
Context *kcontext (Area kstack, void (*entry)(void *), void *arg);

// ----- VME: Virtual Memory -----
bool vme_init (void *(*pgalloc)(int), void (*pgfree)(void *));
void protect (AddrSpace *as);
void unprotect (AddrSpace *as);
void map (AddrSpace *as, void *vaddr, void *paddr, int prot);
Context *ucontext (AddrSpace *as, Area kstack, void *entry);

// ----- MPE: Multi-Processing -----
bool mpe_init (void (*entry)());
int cpu_count (void);
int cpu_current (void);
int atomic_xchg (int *addr, int newval);

#ifdef __cplusplus
}
#endif
```

in abstract-machine/am/include/am.h

PA中有趣的预编译

```
#define _Log(...) \  
do { \  
    printf(__VA_ARGS__); \  
    log_write(__VA_ARGS__); \  
} while (0)
```

为什么要用do{...} while(0) ?

```
#if !defined(likely)  
#define likely(cond)    __builtin_expect(cond, 1)  
#define unlikely(cond)  __builtin_expect(cond, 0)  
#endif
```

likely和unlikely ?

```
#define ANSI_FG_BLACK    "\33[1;30m"  
#define ANSI_FG_RED     "\33[1;31m"  
#define ANSI_FG_GREEN   "\33[1;32m"  
#define ANSI_FG_YELLOW  "\33[1;33m"  
#define ANSI_FG_BLUE    "\33[1;34m"  
#define ANSI_FG_MAGENTA "\33[1;35m"  
#define ANSI_FG_CYAN    "\33[1;36m"  
#define ANSI_FG_WHITE   "\33[1;37m"  
#define ANSI_BG_BLACK   "\33[1;40m"  
#define ANSI_BG_RED     "\33[1;41m"  
#define ANSI_BG_GREEN   "\33[1;42m"  
#define ANSI_BG_YELLOW  "\33[1;43m"  
#define ANSI_BG_BLUE    "\33[1;44m"  
#define ANSI_BG_MAGENTA "\33[1;45m"  
#define ANSI_BG_CYAN    "\33[1;46m"  
#define ANSI_BG_WHITE   "\33[1;47m"  
#define ANSI_NONE       "\33[0m"
```

这个是什么？

PA中有趣的预编译

- 宏维护按键名称列表

```
#define AM_KEYS(_) \
_(ESCAPE) _(F1) _(F2) _(F3) _(F4) _(F5) _(F6) _(F7) _(F8) _(F9) _(F10) _(F11) _(F12) \
_(GRAVE) _(1) _(2) _(3) _(4) _(5) _(6) _(7) _(8) _(9) _(0) _(MINUS) _(EQUALS) _(BACKSPACE) \
_(TAB) _(Q) _(W) _(E) _(R) _(T) _(Y) _(U) _(I) _(O) _(P) _(LEFTBRACKET) _(RIGHTBRACKET) _(BACKSLASH) \
_(CAPSLOCK) _(A) _(S) _(D) _(F) _(G) _(H) _(J) _(K) _(L) _(SEMICOLON) _(APOSTROPHE) _(RETURN) \
_(LSHIFT) _(Z) _(X) _(C) _(V) _(B) _(N) _(M) _(COMMA) _(PERIOD) _(SLASH) _(RSHIFT) \
_(LCTRL) _(APPLICATION) _(LALT) _(SPACE) _(RALT) _(RCTRL) \
_(UP) _(DOWN) _(LEFT) _(RIGHT) _(INSERT) _(DELETE) _(HOME) _(END) _(PAGEUP) _(PAGEDOWN)

#define AM_KEY_NAMES(key) AM_KEY_##key,
enum {
    AM_KEY_NONE = 0,
    AM_KEYS(AM_KEY_NAMES)
};
```



可以使用类似的宏技巧来批量定义变量。以下是几种常见的写法：

1. 批量定义整型变量

```
c

#define DEFINE_INT_VARS(_) \
_(var1) _(var2) _(var3) _(var4) _(var5)

#define DECLARE_INT(var) int var;

// 使用方式
DEFINE_INT_VARS(DECLARE_INT)
// 展开为: int var1; int var2; int var3; int var4; int var5;
```

PA中有趣的预编译

- 阅读PA代码的第一个困难(in nemu/src/nemu-main.c)
 - am_init_monitor()?
 - 还是init_monitor()?

```
int main(int argc, char*argv[]) {  
    /* Initialize the monitor. */  
    #ifdef CONFIG_TARGET_AM  
        am_init_monitor();  
    #else  
        init_monitor(argc, argv);  
    #endif  
    /* Start engine. */  
    engine_start();  
    return is_exit_status_bad();  
}
```

PA1阶段相关的宏

- in nemu/src/monitor/monitor.c

```
#ifndef CONFIG_TARGET_AM
#include <getopt.h>
**
void sdb_set_batch_mode();

static char *log_file = NULL;
static char *diff_so_file = NULL;
static char *img_file = NULL;
static int difftest_port = 1234;

> static long load_img() { ...

> static int parse_args(int argc, char *argv[]) { ...

> void init_monitor(int argc, char *argv[]) { ...
> #else // CONFIG_TARGET_AM...
#endif
```

```
/* Initialize the simple debugger. */
init_sdb();

IFDEF(CONFIG_ITRACE, init_disasm());
```

PA2阶段最重要的宏

- 看看能不能读懂？
 - 可变长宏参数 ()

```
// --- pattern matching wrappers for decode ---
✓ #define INSTPAT(pattern, ...) do { \
    uint64_t key, mask, shift; \
    pattern_decode(pattern, STRLEN(pattern), &key, &mask, &shift); \
✓ if (((uint64_t)INSTPAT_INST(s) >> shift) & mask) == key) { \
    INSTPAT_MATCH(s, ##__VA_ARGS__); \
    goto *(__instpat_end); \
} \
} while (0)

#define INSTPAT_START(name) { const void * __instpat_end = &&concat(__instpat_end_, name);
#define INSTPAT_END(name)    concat(__instpat_end_, name): ; }
```

PA中丰富可用的宏

```
C cpu.h • C macro.h 2 X
ics2025 > nemu > include > C macro.h > MAP(c, f)
16 #ifndef __MACRO_H__
17
18 // calculate the length of an array
19 #define ARRLEN(arr) (int)(sizeof(arr) / sizeof(arr[0]))
20
21 // macro concatenation
22 #define concat_temp(x, y) x ## y
23 #define concat(x, y) concat_temp(x, y)
24 #define concat3(x, y, z) concat(concat(x, y), z)
25 #define concat4(x, y, z, w) concat3(concat(x, y), z, w)
26 #define concat5(x, y, z, v, w) concat4(concat(x, y), z, v, w)
27
28 // macro testing
29 // See https://stackoverflow.com/questions/26684291
30 #define CHOOSE2nd(a, b, ...) b
31 #define MUX_WITH_COMMA(contain_comma, a, b) CHOOSE2nd(contain_comma, a, b)
32 #define MUX_MACRO_PROPERTY(p, macro, a, b) MUX_WITH_COMMA(p, a, b)
33 // define placeholders for some property
34 #define __P_DEF_0 X,
35 #define __P_DEF_1 X,
36 #define __P_ONE_1 X,
37 #define __P_ZERO_0 X,
38 // define some selection functions based on the macro
39 #define MUXDEF(macro, X, Y) MUX_MACRO_PROPERTY(__P_DEF_0, macro, X, Y)
40 #define MUXNDEF(macro, X, Y) MUX_MACRO_PROPERTY(__P_DEF_1, macro, X, Y)
41 #define MUXONE(macro, X, Y) MUX_MACRO_PROPERTY(__P_ONE_1, macro, X, Y)
42 #define MUXZERO(macro, X, Y) MUX_MACRO_PROPERTY(__P_ZERO_0, macro, X, Y)
43
44 // test if a boolean macro is defined
45 #define ISDEF(macro) MUXDEF(macro, 1, 0)
46 // test if a boolean macro is undefined
47 #define ISNDEF(macro) MUXNDEF(macro, 1, 0)
48 // test if a boolean macro is defined to 1
49 #define ISONE(macro) MUXONE(macro, 1, 0)
50 // test if a boolean macro is defined to 0
51 #define ISZERO(macro) MUXZERO(macro, 1, 0)
52
53 // test if a macro of ANY type is defined
54 // NOTE1: it ONLY works inside a function, since it calls `strcmp()`
55 // NOTE2: macros defined to themselves (#define A A) will get wrong results
56 #define isdef(macro) (strcmp("", #macro, "" str(macro)) != 0)
```

```
#define Log(format, ...) \
    _Log(ANSI_FMT("[%s:%d %s] " format, ANSI_FG_BLUE) "\n", \
        __FILE__, __LINE__, __func__, ## __VA_ARGS__)
#define Assert(cond, format, ...) \
    do { \
        if (!(cond)) { \
            MUXDEF(CONFIG_TARGET_AM, printf(ANSI_FMT(format, ANSI_FG_RED) "\n", ## __VA_ARGS__), \
                (fflush(stdout), fprintf(stderr, ANSI_FMT(format, ANSI_FG_RED) "\n", ## __VA_ARGS__)); \
            IFNDEF(CONFIG_TARGET_AM, extern FILE* log_fp; fflush(log_fp)); \
            extern void assert_fail_msg(); \
            assert_fail_msg(); \
            assert(cond); \
        } \
    } while (0)
#define panic(format, ...) Assert(0, format, ## __VA_ARGS__)
#define TODO() panic("please implement me")
```

多路线支持的背后

- ISA_H
 - CFLAGS维护，影响不同路线选择（有兴趣去看看构建过程？）

```
#if defined(CONFIG_ISA_x86)
# define DIFFTEST_REG_SIZE (sizeof(uint32_t) * 9) // GPRs + pc
#elif defined(CONFIG_ISA_mips32)
# define DIFFTEST_REG_SIZE (sizeof(uint32_t) * 38) // GPRs + status + Lo + hi + badvaddr
+ cause + pc
#elif defined(CONFIG_ISA_riscv)
#define RISC_V_GPR_TYPE MUXDEF(CONFIG_RV64, uint64_t, uint32_t)
#define RISC_V_GPR_NUM MUXDEF(CONFIG_RV64, 32, 64)
#define DIFFTEST_REG_SIZE (sizeof(RISC_V_GPR_TYPE) * (RISC_V_GPR_NUM
+ 1))
#elif defined(CONFIG_ISA_loongarch32r)
# define DIFFTEST_REG_SIZE (sizeof(uint32_t) * 33) // GPRs + pc
#else
# error Unsupported ISA
#endif
```

```
#if defined(__ISA_X86__)
# define nemu_trap(code) asm volatile ("int3" : : "a"(code))
#elif defined(__ISA_MIPS32__)
# define nemu_trap(code) asm volatile ("move $v0, %0; sdbbp" : : "a"(code))
#elif defined(__riscv)
# define nemu_trap(code) asm volatile ("mv a0, %0; ebreak" : : "a"(code))
#elif defined(__ISA_LOONGARCH32R__)
# define nemu_trap(code) asm volatile ("move $a0, %0; break 0" : : "a"(code))
#else
# error unsupported ISA __ISA__
#endif
```

```
#if defined(__ARCH_X86_NEMU)
# define DEVICE_BASE 0x0
#else
# define DEVICE_BASE 0xa0000000
#endif

#define MMIO_BASE 0xa0000000

#define SERIAL_PORT (DEVICE_BASE + 0x00003f8)
#define KBD_ADDR (DEVICE_BASE + 0x0000060)
#define RTC_ADDR (DEVICE_BASE + 0x0000048)
#define VGACTL_ADDR (DEVICE_BASE + 0x0000100)
#define AUDIO_ADDR (DEVICE_BASE + 0x0000200)
#define DISK_ADDR (DEVICE_BASE + 0x0000300)
#define FB_ADDR (MMIO_BASE + 0x1000000)
#define AUDIO_SBUF_ADDR (MMIO_BASE + 0x1200000)
```

有趣的预编译

- 发生在实际编译之前
 - 也称为元编程 (meta-programming)
 - Gcc的预处理器同样可以处理汇编代码
 - C++中的模板元编程 ; Rust中的macros ; ...
- Pros
 - 提供灵活的用法 (X-macros)
 - 接近自然语言的写法
- Cons
 - 破坏可读性IOCCC、程序分析 (补全)、.....

```
#define L (  
int main L ) { puts L "Hello, World" ); }
```


编译与链接

(先行剧透本学期的主要内容)

编译、链接

.c → 预编译 → .i → 编译 → .s → 汇编 → .o → 链接 → .out

编译

- 一个不带优化的简易（理想）的编译器
 - C代码中的连续一段总能找到对应的一段连续的机器指令
 - 这就是为什么大家会觉得C是高级的汇编语言！

```
int foo(int n) {  
    int sum = 0;  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum;  
}
```



\$

编译

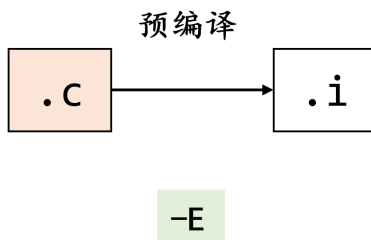
- 一个不带优化的简易（理想）的编译器
 - C代码中的连续一段总能找到对应的一段连续的机器指令
 - 这就是为什么大家会觉得C是高级的汇编语言！

```
int foo(int n) {  
    int sum = 0;  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum;  
}
```

a.c到a.s到a.o

a.s

```
1 foo:
2     n = ARG-1
3     sum = 0
4     i = 1
5     goto .L2
6 .L3:
7     tmp = i
8     sum += tmp
9     i += 1
10 .L2:
11     tmp = i
12     compare (n, tmp)
13     if(<=) goto .L3
14     RETURN-VAL = sum
15
16     ret
```



没有main还不能运行

a.c

```
1 int foo(int n) {
2     int sum = 0;
3     for (int i = 1; i <= n; i++) {
4         sum += i;
5     }
6 }
```

```
$ gcc -c a.c
$ objdump -d a.o
```

a.o: file format elf64-x86-64

Disassembly of section .text:

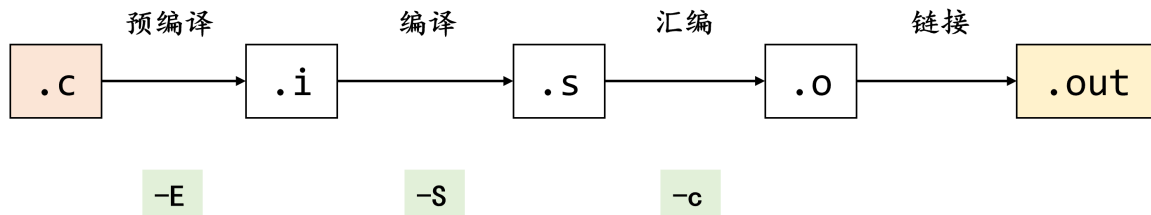
0000000000000000 <foo>:

```
0: 55
1: 48 89 e5
4: 89 7d ec
7: c7 45 fc 00 00 00 00
e: c7 45 f8 01 00 00 00
15: eb 0a
17: 8b 45 f8
1a: 01 45 fc
1d: 83 45 f8 01
21: 8b 45 f8
24: 3b 45 ec
27: 7e ee
29: 8b 45 fc
2c: 5d
2d: c3
```

a.o

```
push    %rbp
mov     %rsp,%rbp
mov     %edi,-0x14(%rbp)
movl    $0x0,-0x4(%rbp)
movl    $0x1,-0x8(%rbp)
jmp     21 <foo+0x21>
mov     -0x8(%rbp),%eax
add     %eax,-0x4(%rbp)
addl    $0x1,-0x8(%rbp)
mov     -0x8(%rbp),%eax
cmp     -0x14(%rbp),%eax
jle     17 <foo+0x17>
mov     -0x4(%rbp),%eax
pop     %rbp
retq
```

链接



- 将多个二进制目标代码拼接在一起
 - C中称为编译单元 (compilation unit)

a.c

```
int foo(int n) {  
    int sum = 0;  
    for (int i = 1; i <= n; i++) {  
        sum += i;  
    }  
    return sum;  
}
```

gcc -c a.c

a.o

b.c

```
#include <stdio.h>  
int foo(int n);  
int main(){  
    printf("%d\n"),foo(100));  
}
```

gcc -c b.c

b.o

链接

?



\$

链接

- 将多个二进制目标代码拼接在一起
 - C中称为编译单元 (compilation unit)
 - 甚至可以链接C++ , rust , ...代码

```
extern "C" {  
    int foo() { return 0; }  
}
```

```
int bar() { return 0; }
```

```
$ g++ -c a.cc  
$ objdump -d a.o
```

a.o: file format elf64-x86-64

Disassembly of section .text:

0000000000000000 <foo>:

0:	55	push	%rbp
1:	48 89 e5	mov	%rsp,%rbp
4:	b8 00 00 00 00	mov	\$0x0,%eax
9:	5d	pop	%rbp
a:	c3	retq	

000000000000000b <_Z3barv>:

b:	55	push	%rbp
c:	48 89 e5	mov	%rsp,%rbp
f:	b8 00 00 00 00	mov	\$0x0,%eax
14:	5d	pop	%rbp
15:	c3	retq	

PA中的细节

```
#ifdef cplusplus
extern "C" {
#endif

// ----- TRM: Turing Machine -----
extern Area heap;
void putch (char ch);
void halt (int code) __attribute__((__noreturn__));

// ----- IOE: Input/Output Devices -----
bool ioe_init (void);
void ioe_read (int reg, void *buf);
void ioe_write (int reg, void *buf);
#include "amdev.h"

// ----- CTE: Interrupt Handling and Context Switching -----
bool cte_init (Context *(*handler)(Event ev, Context *ctx));
void yield (void);
bool ienabled (void);
void iset (bool enable);
Context *kcontext (Area kstack, void (*entry)(void *), void *arg);

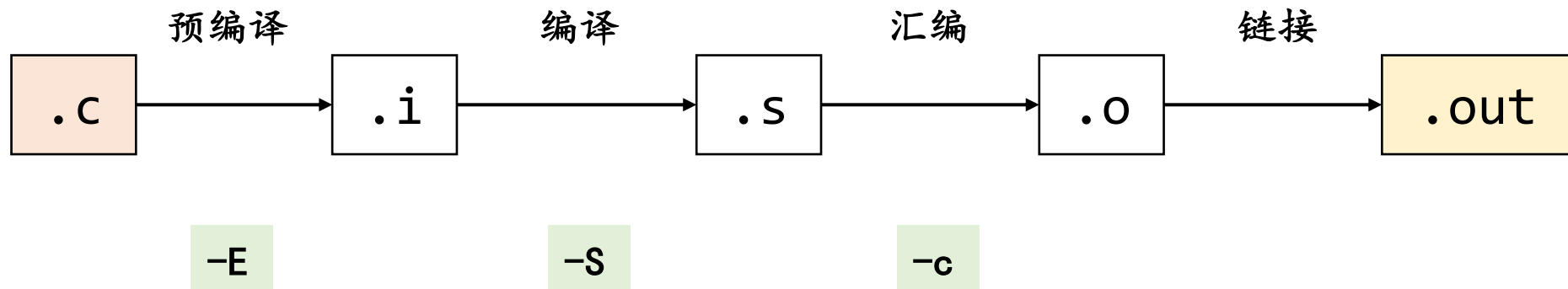
// ----- VME: Virtual Memory -----
bool vme_init (void *(*pgalloc)(int), void (*pgfree)(void *));
void protect (AddrSpace *as);
void unprotect (AddrSpace *as);
void map (AddrSpace *as, void *vaddr, void *paddr, int prot);
Context *ucontext (AddrSpace *as, Area kstack, void *entry);

// ----- MPE: Multi-Processing -----
bool mpe_init (void (*entry)());
int cpu_count (void);
int cpu_current (void);
int atomic_xchg (int *addr, int newval);

#ifdef __cplusplus
}
#endif
```

in abstract-machine/am/include/am.h

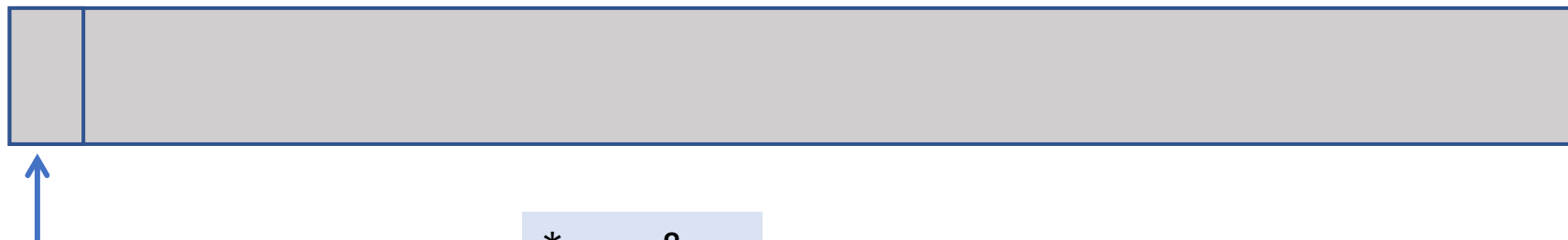
从源代码到可执行文件



加载：进入C语言的世界

C程序执行的两个视角

- 静态：C 代码的连续一段总能对应到一段连续的机器指令
- 动态：C 代码执行的状态总能对应到机器的状态
 - 源代码视角
 - 函数、变量、指针.....
 - 机器指令视角
 - 寄存器、内存、地址.....
- 两个视角的共同之处：内存
 - 代码、变量 (源代码视角) = 地址 + 长度 (机器指令视角)
 - (不太严谨地) 内存 = 代码 + 数据 + 堆栈
 - 因此理解 C 程序执行最重要的就是内存模型



$x=1 \rightarrow$

<pre>*p = &x *p = 1</pre>

```
int main(int argc, char *argv[]) {  
    int *p = (void *) 1;    //OK  
    *p = 1;    //Segmentation fault  
}
```

内存只是个**字节序列**

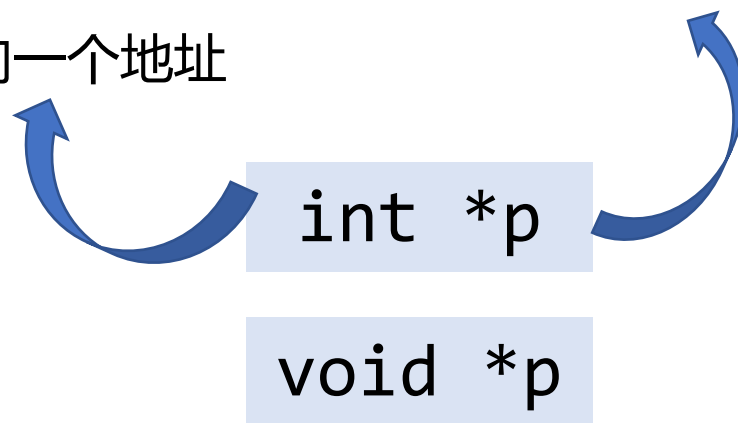
无论何种类型的指针都只是**地址** + 对指向内存的**解读**

内存



对指向的内存的解读：**int**

指针：指向一个地址



内存

对a地址开始的内存的解读：**long**



8个字节，64位bit

a

long *

取此块内存的值

*** (long *)**

p指针：指向一个地址a

%p输出地址a

void *p

```
void printptr(void *p) {  
    printf("p = %p; *p = %016lx\n", p, *(long *)p);  
}
```

指向的内存解读为long输出

输出指针指向的地址

以16位16进制数格式输出：16*4=64bit

main, argc和argv

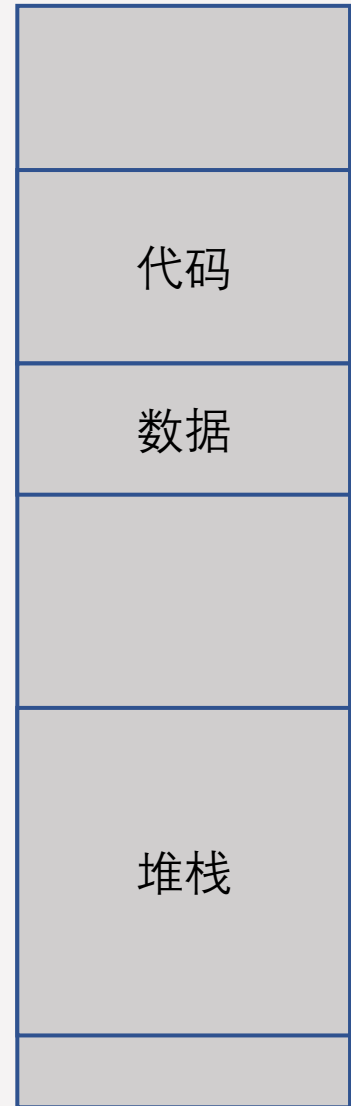
- 一切皆可取地址！

```
void printptr(void *p) {  
    printf("p = %p; *p = %016lx\n", p, *(long *)p);  
}  
int x;  
int main(int argc, char *argv[]) {  
    printptr(main); // 代码  
    printptr(&main);  
    printptr(&x); // 数据  
    printptr(&argc); // 堆栈  
    printptr(argv);  
    printptr(&argv);  
    printptr(argv[0]);  
}
```

指向的地址处内存解读为long输出16位16进制

输出指针指向的地址

\$



C Type System

- **类型**：对一段内存的**解读方式**
 - 非常汇编：没有class, polymorphism, type traits,
 - C里的所有的数据都可以理解成**地址（指针）+类型（对地址的解读）**
- 例子（是不是感到学到了假了C语言）

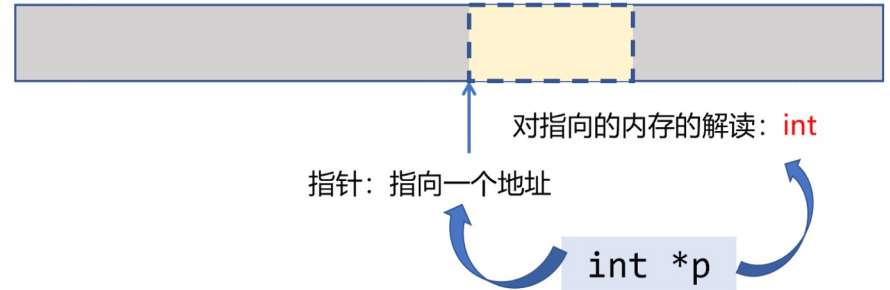
```
int main(int argc, char *argv[]) {  
    int (*f)(int, char *[]) = main;  
    if (argc != 0) {  
        char ***a = &argv, *first = argv[0], ch = argv[0][0];  
        printf("arg = \"%s\"; ch = '%c'\n", first, ch);  
        assert(**a == ch);  
        f(argc - 1, argv + 1);  
    }  
}
```

```
$ ./a.out 1 2 3 hello  
arg = "./a.out"; ch = '.'  
arg = "1"; ch = '1'  
arg = "2"; ch = '2'  
arg = "3"; ch = '3'  
arg = "hello"; ch = 'h'
```



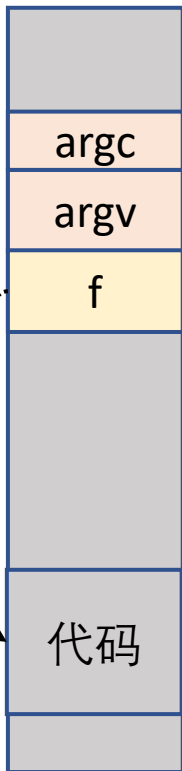
```
int main(int argc, char *argv[]) {  
    int (*f)(int, char *[]) = main;  
    if (argc != 0) {  
        char ***a = &argv, *first = argv[0], ch = argv[0][0];  
        printf("arg = \"%s\""; ch  
        assert(**a == ch);  
        f(argc - 1, argv + 1);  
    }  
}
```

函数指针f



内存

堆栈



char *argv[] → char **argv → (char *)*argv

4字节

? 指针, 存储地址, 64位机器, 8字节

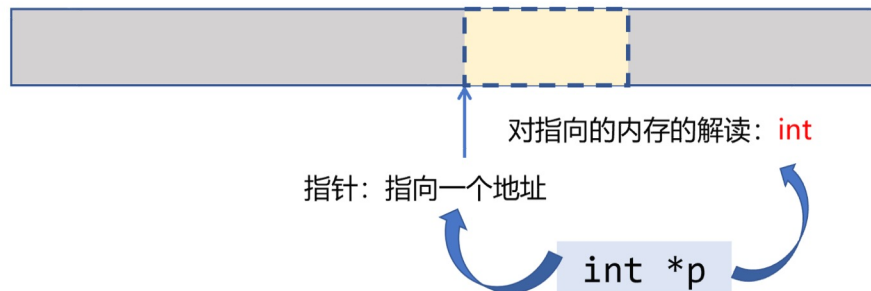
? 指针, 存储地址, 64位机器, 8字节

main

代码



```
int main(int argc, char *argv[]) {  
    int (*f)(int, char *[]) = main;  
    if (argc != 0) {  
        char ***a = &argv, *first  
        printf("arg = \"%s\""; ch  
        assert(**a == ch);  
        f(argc - 1, argv + 1);  
    }  
}
```

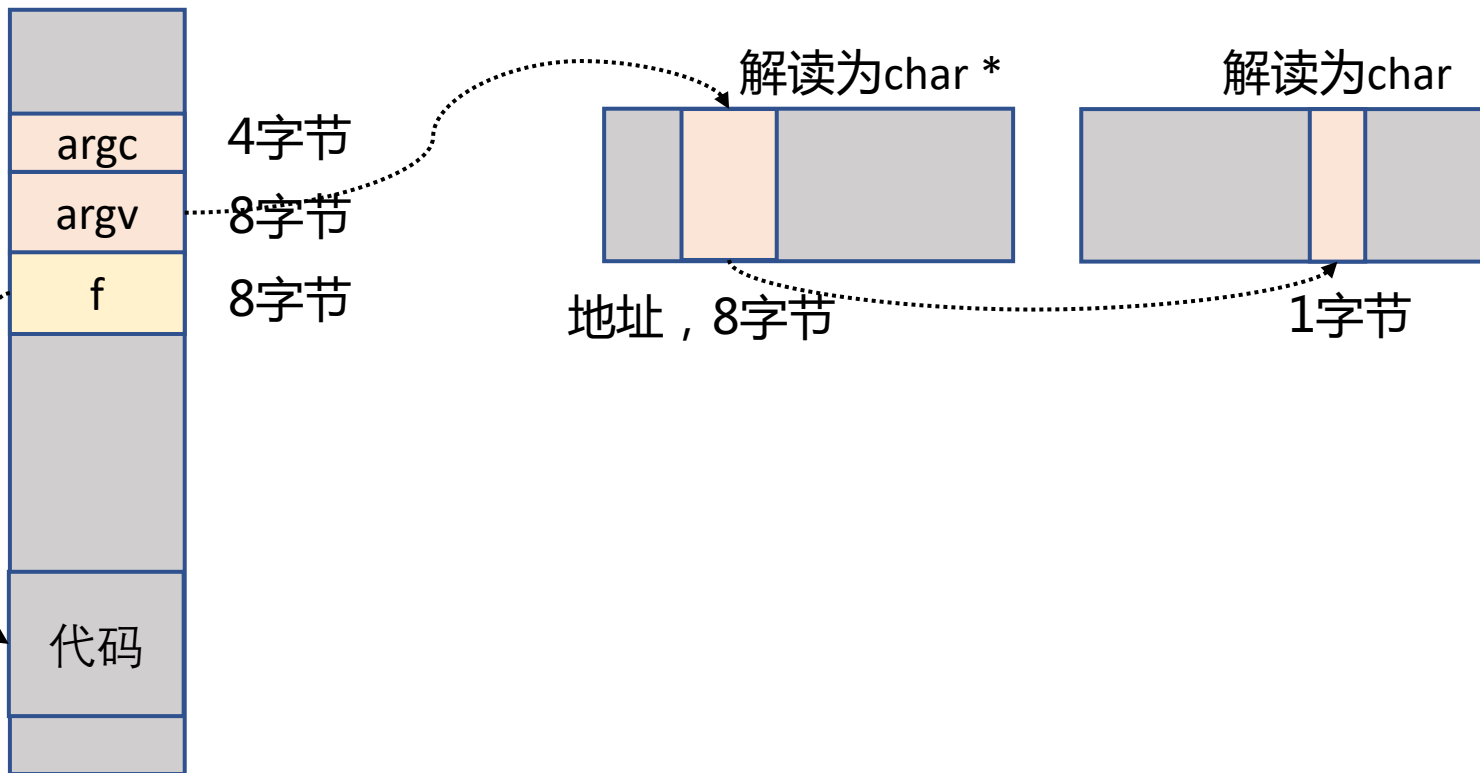


char *argv[] → char **argv → (char *)*argv

内存

堆栈

main



```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0], ch = argv[0][0];
        printf("arg = \"%s\"; ch = '%c'\n", first, ch);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

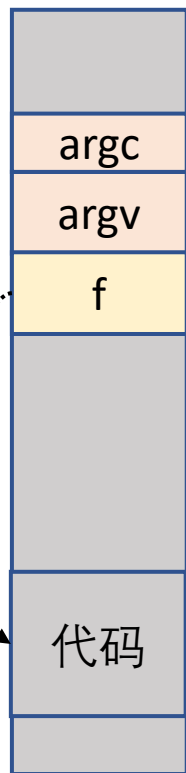
```



内存

堆栈

main



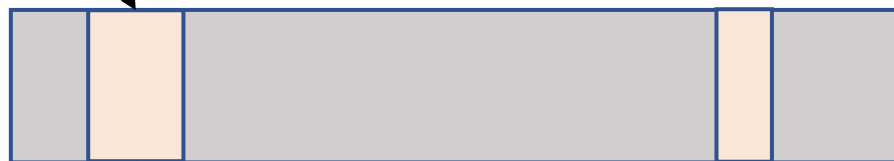
4字节

8字节

8字节

代码

解读为char *



地址, 8字节

1字节

解读为char

```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0], ch = argv[0][0];
        printf("arg = \"%s\"; ch = '%c'\n", first, ch);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```

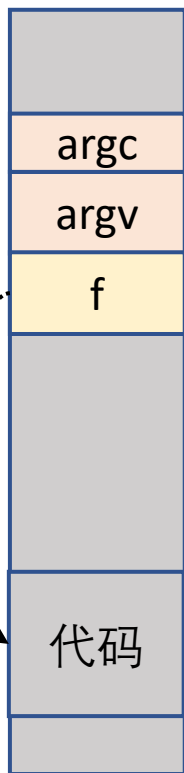


什么是argv+1 ?

内存

堆栈

main



4字节

8字节

8字节

解读为char *

地址, 8字节

1字节

解读为char

int *p; → p+1
按int解读下一个内存地址

```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0], ch = argv[0][0];
        printf("arg = \"%s\"; ch = '%c'\n", first, ch);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

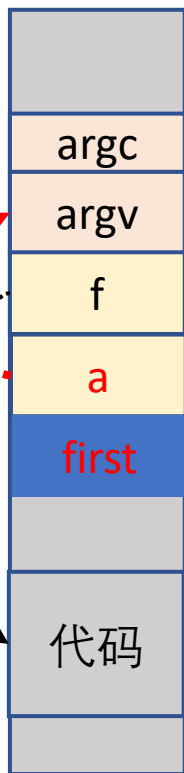
```



内存

堆栈

main



4字节

8字节

8字节

8字节

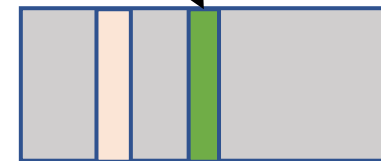
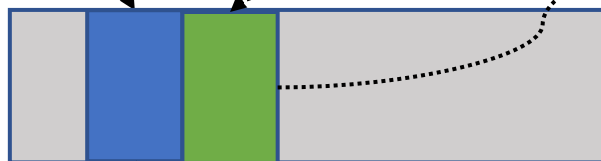
8字节

地址, 8字节

按char解读

argv+1、
&argv[1]

解读为char *



1字节

解读为char

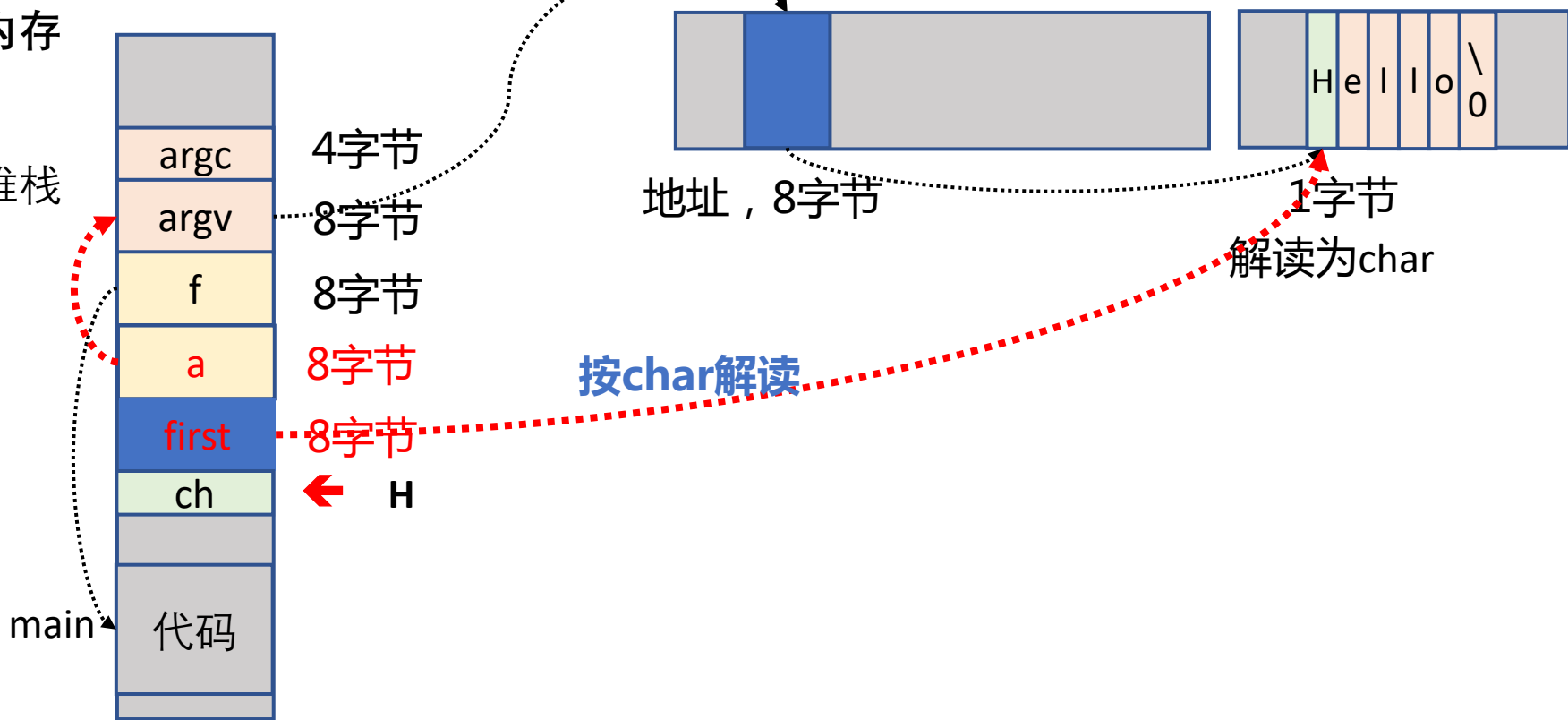
```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0], ch = argv[0][0];
        printf("arg = \"%s\"; ch = '%c'\n", first, ch);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```

内存

堆栈




```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0];
        printf("arg = \"%s\"; ch = '%c'\\n", *a, *first);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```

↓

```

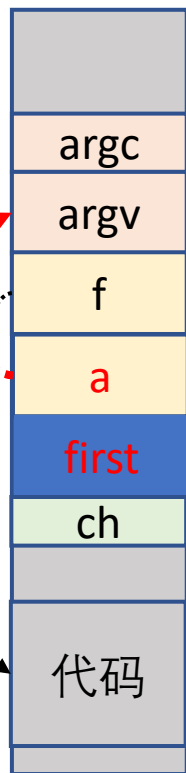
$ ./a.out 1 2 3 hello
arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'
arg = "hello"; ch = 'h'

```

内存

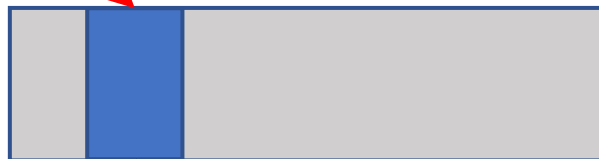
堆栈

main



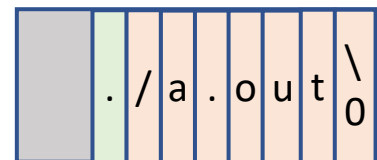
5

解读为char *



地址, 8字节

Hello



1字节

解读为char

按char解读

每一轮输出argv指向的地址所存地址处内存的字符串和第一个字符

arg = "./a.out"; ch = '.'

./a.out 1 2 3 Hello

```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0];
        printf("arg = \"%s\"; ch = '%c'\\n", *a, **first);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```



```

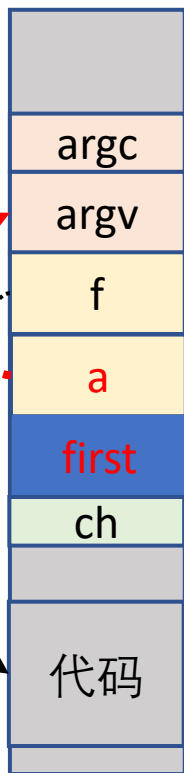
$ ./a.out 1 2 3 hello
arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'
arg = "hello"; ch = 'h'

```

内存

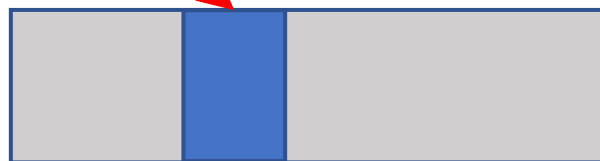
堆栈

main



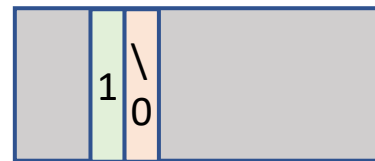
4

解读为char *



地址, 8字节

Hello



1字节

解读为char

按char解读

每一轮输出argv指向的地址所存地址处内存的字符串和第一个字符

```

arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'

```

./a.out 1 2 3 Hello

```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0];
        printf("arg = \"%s\"; ch = '%c'\\n", *a, *first);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```

↓

```

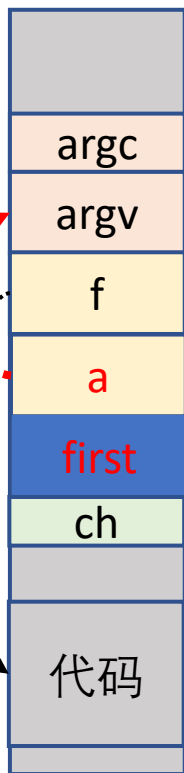
$ ./a.out 1 2 3 hello
arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'
arg = "hello"; ch = 'h'

```

内存

堆栈

main



3

地址, 8字节

按char解读

解读为char *

Hello

1字节

解读为char

每一轮输出argv指向的地址所存地址处内存的字符串和第一个字符

arg = "./a.out"; ch = '.'

arg = "1"; ch = '1'

arg = "2"; ch = '2'

./a.out 1 2 3 Hello

```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0];
        printf("arg = \"%s\"; ch = '%c'\\n", *first, **a);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```

```

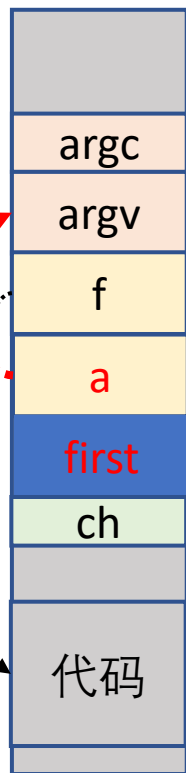
$ ./a.out 1 2 3 hello
arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'
arg = "hello"; ch = 'h'

```

内存

堆栈

main



2

地址, 8字节

按char解读

解读为char*

Hello

1字节

解读为char

每一轮输出argv指向的地址所存地址处内存的字符串和第一个字符

```

arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'

```

./a.out 1 2 3 Hello

```

int main(int argc, char *argv[]) {
    int (*f)(int, char *[]) = main;
    if (argc != 0) {
        char ***a = &argv, *first = argv[0];
        printf("arg = \"%s\"; ch = '%c'\\n", *first, **a);
        assert(**a == ch);
        f(argc - 1, argv + 1);
    }
}

```

↓

```

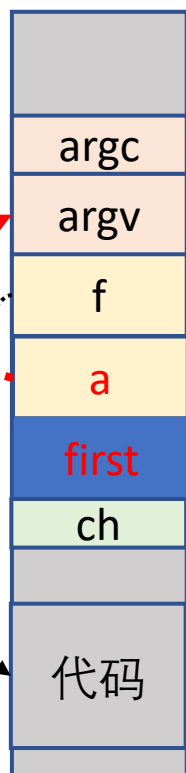
$ ./a.out 1 2 3 hello
arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'
arg = "hello"; ch = 'h'

```

内存

堆栈

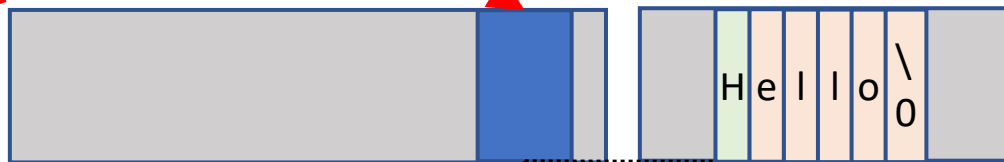
main



1

地址, 8字节

按char解读



Hello

1字节

解读为char

每一轮输出argv指向的地址所存地址处内存的字符串和第一个字符

arg = "./a.out"; ch = '.'
arg = "1"; ch = '1'
arg = "2"; ch = '2'
arg = "3"; ch = '3'
arg = "hello"; ch = 'h'

\$

```
void printptr(void *p) { 指向的地址处内存解读为long输出16位16进制
    printf("p = %p; *p = %016lx\n", p, *(long *)p);
}
```

输出指针指向的地址

```
int x;
int main(int argc, char *argv[]) {
    printptr(main); // 代码
    printptr(&main);
    printptr(&x); // 数据
    printptr(&argc); // 堆栈
    printptr(argv);
    printptr(&argv);
    printptr(argv[0]);
}
```

```
$ ./a.out
```

```
p = 0x563afd3cf163; *p = 10ec8348e5894855
p = 0x563afd3cf163; *
p = 0x563afd3d2034; *
p = 0x7ffcada0761c; *p = fd3cf1d000000001
p = 0x7ffcada07708; *p = 00007ffcada084fe
p = 0x7ffcada07610; *p = 00007ffcada07708
p = 0x7ffcada084fe; *p = 0074756f2e612f2e
```

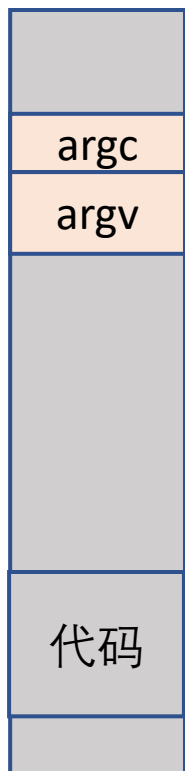
地址XX, 地址XX内存的内容

内存

堆栈

a→

b→



4字节

8字节

地址, 8字节

1字节

解读为char

	p	*p
printptr(&argc)	a	argc值
printptr(argv)	c	d
printptr(&argv)	b	c
printptr(argv[0])	d	d处字符串

从main函数开始执行

- 标准规定C程序从main开始执行
 - （思考题：谁调用的main？进程执行的第一条指令是什么？）

```
int main(int argc, char *argv[]);
```

- argc (argument count): 参数个数
- argv (argument vector): 参数列表（NULL结束）
- `ls -al`
 - `argc = 2, argv = ["ls", "-al", NULL]`

核心准则：编写可读代码

怎样写代码才能从一个大型项目中存活下来？

核心准则：编写可读代码

一个极端不可读的例子

- [IOCCC'11 best self documenting program](#)
 - 不可读 = 不可维护

```
puts(usage: calculator 11/26+222/31
+~~~~~calculator-\
!              7.584,367 )
+~~~~~+
! clear ! 0 ||1 -x 1 tan I (/) |
+~~~~~+
! 1 | 2 | 3 ||1 1/x 1 cos I (*) |
+~~~~~+
! 4 | 5 | 6 ||1 exp 1 sqrt I (+) |
+~~~~~+
! 7 | 8 | 9 ||1 sin 1 log I (-) |
+~~~~~(0 )
```

一个极端不可读的例子

- [IOCCC'11 best self documenting program](#)
 - 不可读 = 不可维护

```
#define clear 1;
    if(c>=11){c=0;sscanf(_, "%lf%c",&r,&c);while(*++_-
c);} \    else if(argc>=4&&!main(4-
(*_++=='('),argv))_++;g:c+=
#define puts(d,e) return 0;}{double a;int b;char
    c=(argc<4?d)&15;\ b=(*_%__LINE__+7)%9*(3*e>>c&1);c+=
#define I(d)
(r);if(argc<4&&*#d==*_){a=r;r=usage?r*a:r+a;goto
g;}c=c
#define return if(argc==2)printf("%f\n",r);return
argc>=4+ #define usage main(4-__LINE__/26,argv)
#define calculator *_*(int)
#define l (r);r=--b?r:
#define _ argv[1]
```

\$ █

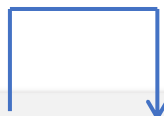
一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例

```
void (*signal (int sig, void (*func)(int)))(int);
```

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例

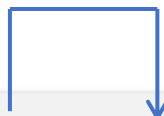


```
void (*signal (int sig, void (*func)(int)))(int);
```

- signal是一个函数
 - 参数为.....
 - 返回值为.....

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例



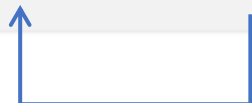
```
void (*signal (int sig, void (*func)(int)))(int);
```

- signal是一个函数
 - 参数为int和一个函数指针 (参数为int , 返回值为void)
 - 返回值为.....

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例

```
void (*signal (int sig, void (*func)(int)))(int);
```



- signal是一个函数
 - 参数为int和一个函数指针 (参数为int , 返回值为void)
 - 返回值为一个.....指针

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例



```
void (*signal (int sig, void (*func)(int))) (int);
```

- signal是一个函数
 - 参数为int和一个指向函数的指针 (函数参数为int , 返回值为void)
 - 返回值为一个指向函数的指针 (参数为... , 返回值为...)

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例



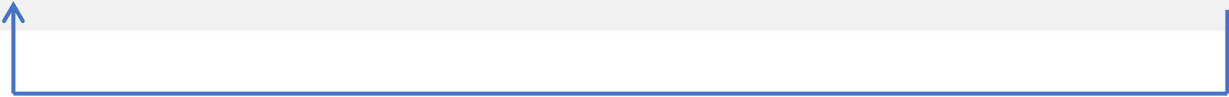
```
void (*signal (int sig, void (*func)(int))) (int);
```

- signal是一个函数
 - 参数为int和一个指向函数的指针 (函数参数为int , 返回值为void)
 - 返回值为一个指向函数的指针 (参数为int , 返回值为...)

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则的案例

```
void (*signal (int sig, void (*func)(int)))(int);
```



- signal是一个函数
 - 参数为int和一个指向函数的指针 (函数参数为int , 返回值为void)
 - 返回值为一个指向函数的指针 (参数为int , 返回值为void)
- 太复杂了>_< , 有没有更简单的办法

一个现实中可能遇到的例子

- 人类不可读版本 (STFW: clockwise/spiral rule)

- 终极使用顺时针螺旋法则

```
void (*signal (int sig, void (*func)(int)))(int);
```

- signal是一个函数
 - 参数为int和一个指向函数的指针 (函数参数为int , 返回值为void)
 - 返回值为一个指向函数的指针 (参数为int , 返回值为void)

```
void (*signal (int sig, void (*func)(int)))(int);
```

```
void (*signal (int sig, func))(int);
```

```
void (*)(int);
```

一个现实中可能遇到的例子

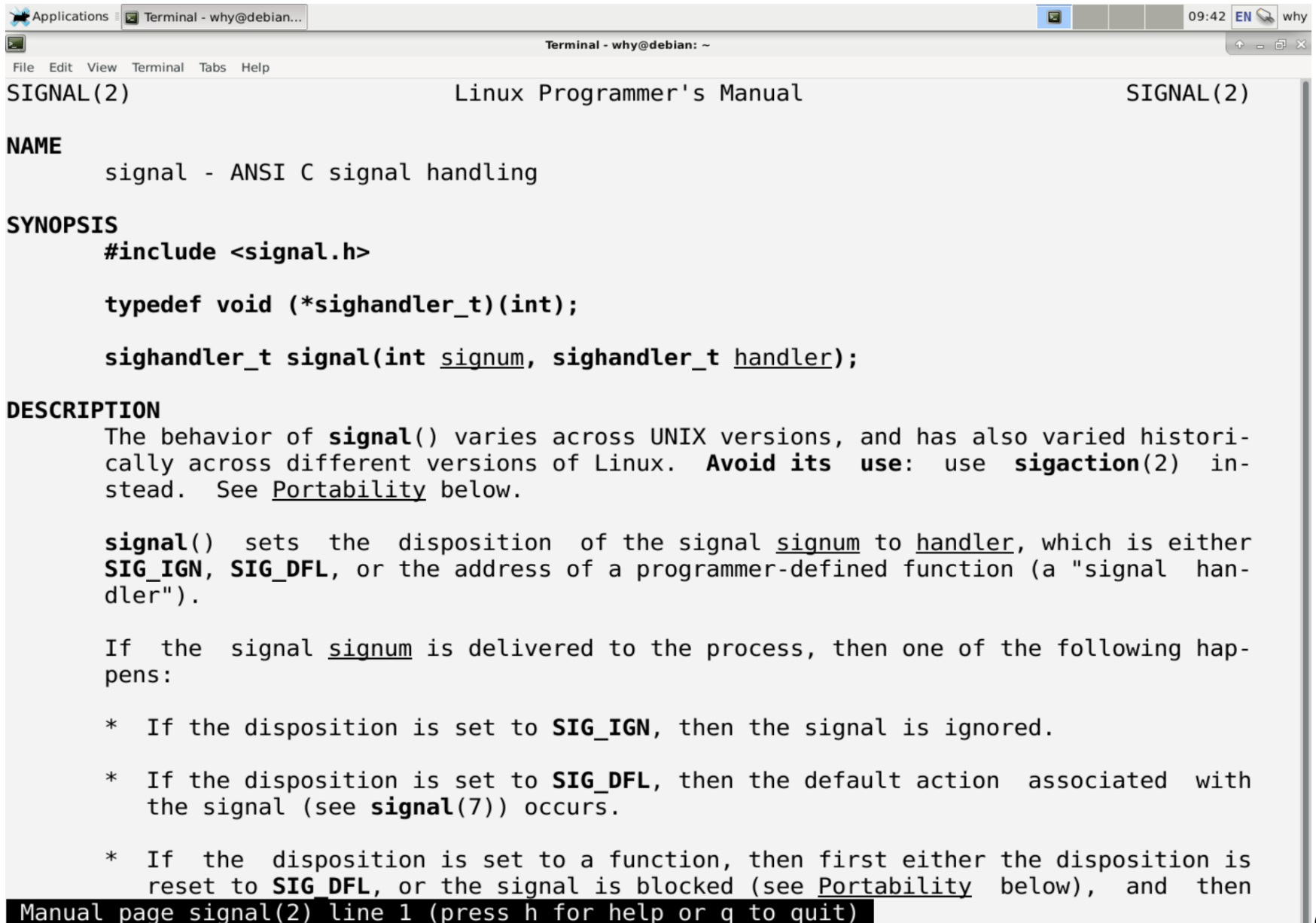
- 人类不可读版本 (STFW: clockwise/spiral rule)
 - 终极使用顺时针螺旋法则

```
void (*signal (int sig, void (*func)(int)))(int);
```

- 人类可读版本

```
typedef void (*sighandler_t)(int);  
sighandler_t signal(int, sighandler_t);
```

man 2 signal



```
Applications Terminal - why@debian... 09:42 EN why
Terminal - why@debian: ~
File Edit View Terminal Tabs Help
SIGNAL(2) Linux Programmer's Manual SIGNAL(2)

NAME
    signal - ANSI C signal handling

SYNOPSIS
    #include <signal.h>

    typedef void (*sighandler_t)(int);

    sighandler_t signal(int signum, sighandler_t handler);

DESCRIPTION
    The behavior of signal() varies across UNIX versions, and has also varied historically across different versions of Linux. Avoid its use: use sigaction(2) instead. See Portability below.

    signal() sets the disposition of the signal signum to handler, which is either SIG_IGN, SIG_DFL, or the address of a programmer-defined function (a "signal handler").

    If the signal signum is delivered to the process, then one of the following happens:

    * If the disposition is set to SIG_IGN, then the signal is ignored.

    * If the disposition is set to SIG_DFL, then the default action associated with the signal (see signal(7)) occurs.

    * If the disposition is set to a function, then first either the disposition is reset to SIG_DFL, or the signal is blocked (see Portability below), and then

Manual page signal(2) line 1 (press h for help or q to quit)
```

编写代码的准则：降低维护成本

Programs are meant to be read by humans and only incidentally for computers to execute. — D. E. Knuth

(程序首先是拿给人读的，其次才是被机器执行。)

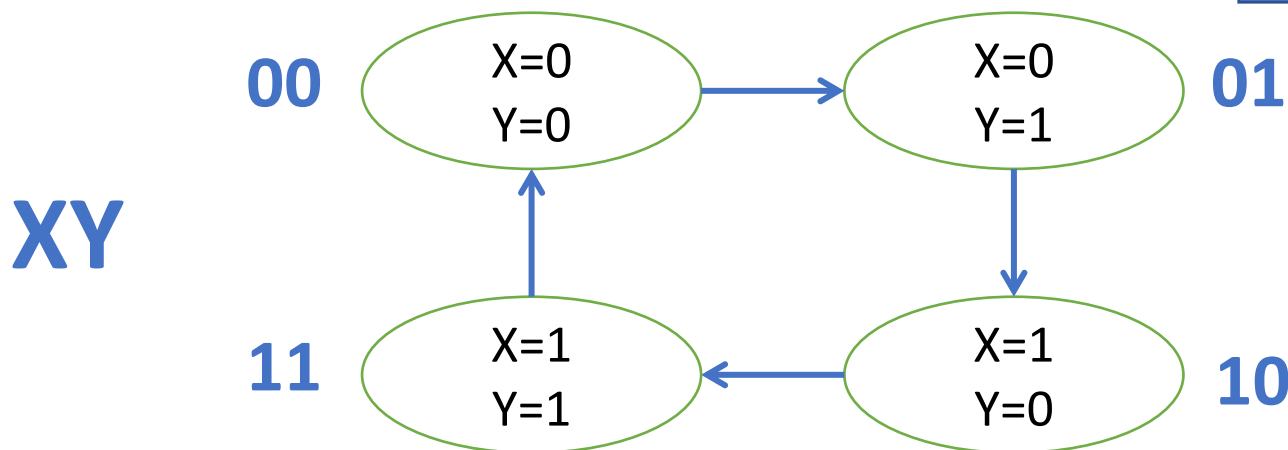
- 宏观
 - 做好分解和解耦（现实世界也是这样管理复杂系统）
- 微观
 - “不言自明”
 - 通过阅读代码能够理解一段程序的行为（implementation）
 - “不言自证”
 - 通过阅读代码能够验证一段程序implementation与specification的一致性

例子：实现数字逻辑电路模拟器

数字逻辑电路模拟器

- 假想的数字逻辑电路

- 若干个1-bit边沿触发寄存器 (X, Y,)
- 若干个逻辑门



- 基本思路：状态（存储模拟）+计算模拟

- 状态 = 变量
 - int X = 0, Y = 0;
- 计算

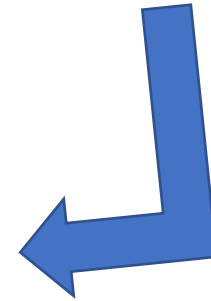
```
X1 = (!X && Y) || (X && !Y);  
Y1 = !Y;  
X = X1; Y = Y1;
```

数字逻辑电路模拟器

- 需求
 - 加一位边沿寄存器
 - 自己独立的逻辑

```
int X=0, Y=0;
int X1=0, Y1=0;
while(1){
    X1 = (!X&&Y) || (X&&!Y);
    Y1 = !Y;
    X = X1; Y = Y1;
}
```

```
int X=0, Y=0, Z=0;
int X1=0, Y1=0, Z1=0;
while(1){
    X1 = (!X&&Y) || (X&&!Y);
    Y1 = !Y;
    Z1 = .....;
    X = X1; Y = Y1; Z = Z1;
}
```



通用数字逻辑电路模拟器

```
#define FORALL_REGS(_)    _(X) _(Y)
#define LOGIC             X1 = (!X&&Y)|| (X&&!Y); \
                          Y1 = !Y;

#define DEFINE(X)         static int X, X##1;
#define UPDATE(X)         X = X##1;
#define PRINT(X)          printf("#X " = %d; ", X);
int main() {
    FORALL_REGS(DEFINE);

    while (1) { // clock
        FORALL_REGS(PRINT);
        putchar('\n');
        sleep(1);
        LOGIC;
        FORALL_REGS(UPDATE);
    }
}
```

File Edit View Terminal Tabs Help

```
1 #include <stdio.h>
2 #include <unistd.h>
3 #define FORALL_REGS(_) _(X) _(Y)
4 #define LOGIC          X1 = (!X && Y) || (X&&!Y); \
5                        Y1 = !Y;
6 #define DEFINE(X)      static int X, X##1;
7 #define UPDATE(X)      X = X##1;
8 #define PRINT(X)       printf(#X " = %d; ", X);
9
10 int main() {
11     FORALL_REGS(DEFINE);
12     while (1) { // clock
13         FORALL_REGS(PRINT); putchar('\n'); sleep(1);
14         LOGIC;
15         FORALL_REGS(UPDATE);
16     }
17 }
18
```

a.c 3,1 All
"a.c" 18L, 437C

Terminal - a.c (~/ICS_teach) - VIM

File Edit View Terminal Tabs Help

```
1 #include <stdio.h>
2 #include <unistd.h>
3 #define FORALL_REGS(_)(X)_(Y)_(Z)
4 #define LOGIC          X1 = (!X&&Y&&Z)|| (X&&(!(Y&&Z))); \
5                        Y1 = (!Y&&Z)|| (Y&&!Z); \
6                        Z1 = !Z;
7 #define DEFINE(X)      static int X, X##1;
8 #define UPDATE(X)      X = X##1;
9 #define PRINT(X)       printf(#X " = %d; ", X);
10
11 int main() {
12     FORALL_REGS(DEFINE);
13     while (1) { // clock
14         FORALL_REGS(PRINT); putchar('\n'); sleep(1);
15         LOGIC;
16         FORALL_REGS(UPDATE);
17     }
18 }
19
```

a.c4,1All
"a.c" 19L, 498C

通用数字逻辑电路模拟器

```
#define FORALL_REGS(_)    _(X) _(Y)
#define LOGIC             X1 = (!X&&Y)|| (X&&!Y); \
                          Y1 = !Y;

#define DEFINE(X)         static int X, X##1;
#define UPDATE(X)         X = X##1;
#define PRINT(X)          printf("#X " = %d; ", X);
int main() {
    FORALL_REGS(DEFINE);

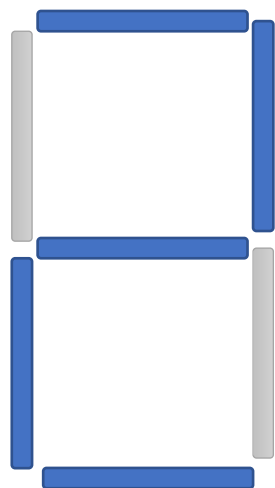
    while (1) { // clock
        FORALL_REGS(PRINT);
        putchar('\n');
        sleep(1);
        LOGIC;
        FORALL_REGS(UPDATE);
    }
}
```

使用预编译：Pros and Cons

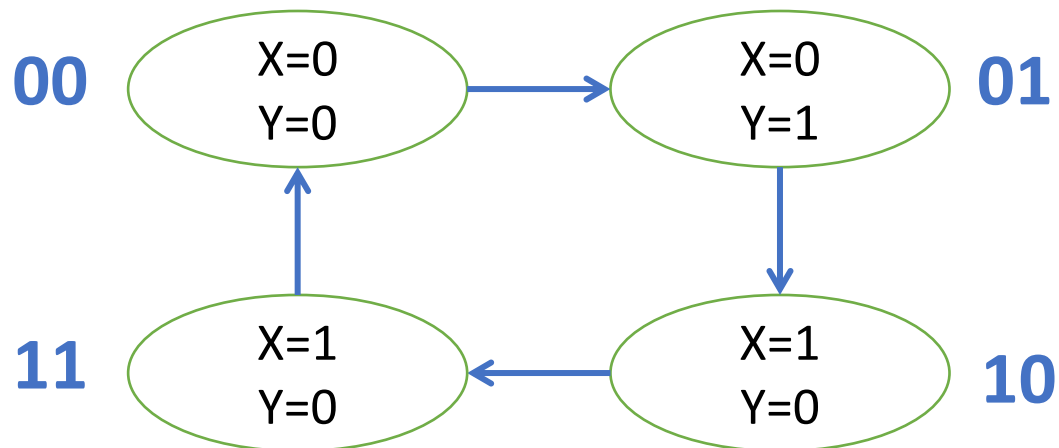
- Pros
 - 增加/删除寄存器只要改很少的地方
 - 阻止一些编程错误
 - 忘记更新寄存器
 - 忘记打印寄存器
 - “不言自明”
- Cons
 - 可读性变差（不太像C代码）
 - “不言自证” 差一些
 - 给IDE解析带来一些困难

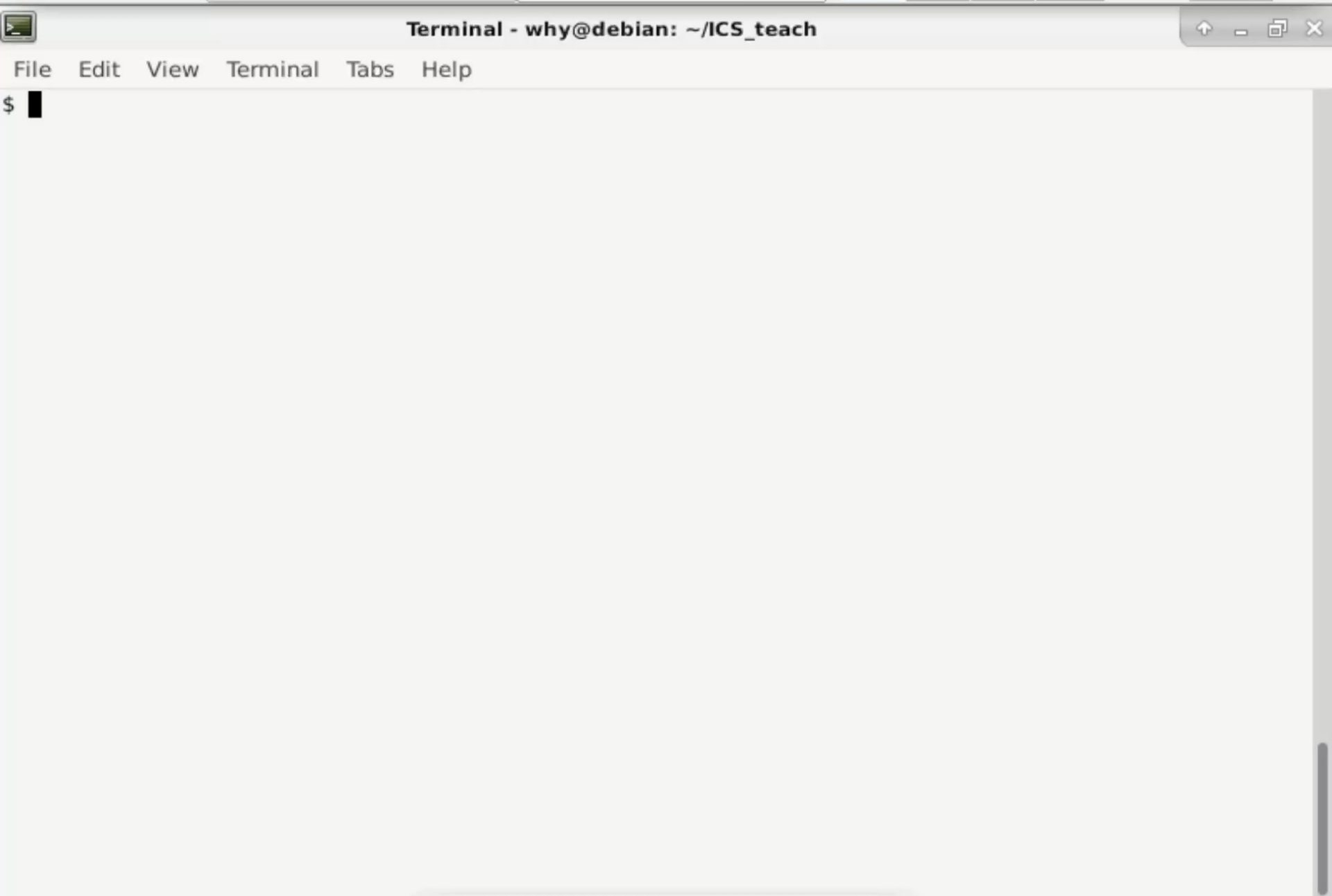
更完整的实现：数码管显示

- logisim.c 和 display.py



XY





```

#include <stdio.h>
#include <unistd.h>
#define FORALL_REGS(_)      _(X) _(Y)
#define OUTPUTS(_)         _(A) _(B) _(C) _(D) _(E) _(F) _(G)
#define LOGIC
                            X1 = (!X && Y) || (X && !Y); \
                            Y1 = !Y; \
                            A  = (!X && !Y) || (X && !Y) || (X && Y); \
                            B  = 1; \
                            C  = (!X && !Y) || (!X && Y) || (X && Y); \
                            D  = (!X && !Y) || (X && !Y) || (X && Y); \
                            E  = (!X && !Y) || (X && !Y); \
                            F  = (!X && !Y); \
                            G  = (X && !Y) || (X && Y);
#define DEFINE(X)           static int X, X##1;
#define UPDATE(X)          X = X##1;
#define PRINT(X)            printf(#X " = %d; ", X);
int main() {
    FORALL_REGS(DEFINE);
    OUTPUTS(DEFINE);
    while (1) { // clock
        LOGIC;
        OUTPUTS(PRINT);
        putchar('\n');
        fflush(stdout);
        sleep(1);
        FORALL_REGS(UPDATE);
    }
}

```

更完整的实现：数码管显示

- logisim.c 和 display.py
 - 也可以考虑增加更多外设：开关、UART等
 - 原理无限接近数字电路课玩过的FPGA
- 等等.....FPGA？
 - 这玩意不是万能的吗？？？
 - 我们能模拟它，是不是就能模拟计算机系统？
 - Yes!
 - 我们实现了一个超级超级低配版 NEMU!



例子：实现YEMU全系统模拟器

更多的计算机系统模拟器

- am-kernels/litenes
 - 一个 “最小” 的 NES 模拟器
 - 自带硬编码的 ROM 文件
- fceux-am
 - 一个非常完整的高性能 NES 模拟器
 - 包含对卡带定制芯片的模拟 (src/boards)
- QEMU
 - 工业级的全系统模拟器
 - 2011 年发布 1.0 版本
 - 有兴趣的同学可以 RTFSC
 - 作者：传奇黑客 Fabrice Bellard

永远不要停止对好代码的追求

End.

- C语言简单（在可控时间成本里可以精通）
- C语言通用（大量系统是C语言编写的）
- C语言实现对底层机器的精准控制（鸿蒙）
- 推荐阅读：[The Art of Readable Code](#)