"造轮子"的方法和乐趣

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本讲概述

《计算机系统基础》课到底学了啥?

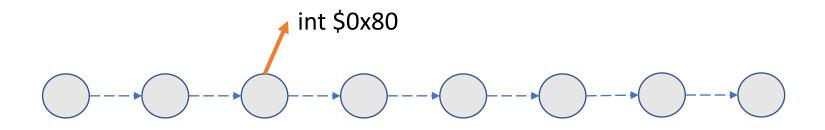
- 本讲内容
 - 一个关于编译运行的问题
 - 那些年我们造过的轮子

一个关于编译运行的问题

回到第一次C语言课

- 现在, 一位同学对这个过程提出了质疑
 - 我不信! 我就觉得是 gcc 一个程序直接搞定的
 - 道理上完全可以这么实现
 - 如何说服这位同学?

所有程序的执行都是状态机



```
$ ls -l
total 4
-rw-rw-r-- 1 why why 14 12月 16 08:19 a.c
$ gcc a.c
$ ls -l
total 20
-rw-rw-r-- 1 why why 14 12月 16 08:19 a.c
-rwxrwxr-x 1 why why 15912 12月 16 08:19 a.out
```

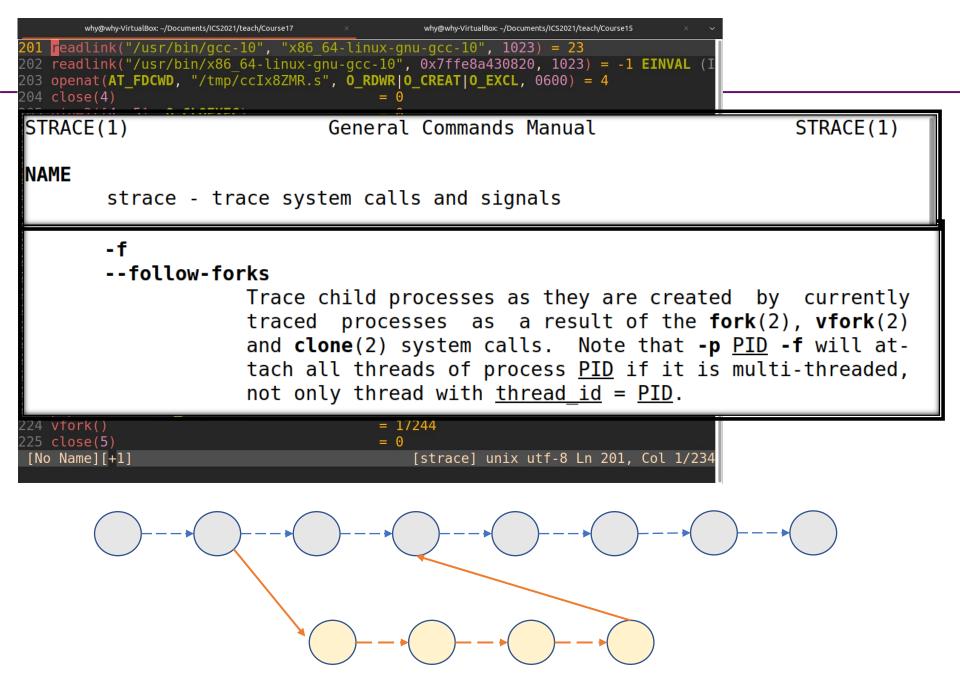
答案: 用工具!

- 观察 trace: 查看 gcc 是否调用了其他命令
 - strace -f -qq gcc a.c 2>&1 | vim -
 - 可以使用 grep (shell) 或 :g! (vim) 筛选感兴趣的系统调用

```
STRACE(1)
                          General Commands Manual
                                                                   STRACE(1)
NAME
       strace - trace system calls and signals
SYNOPSIS
       strace [-ACdffhikqqrtttTvVwxxyyzZ] [-I n] [-b execve] [-e expr]...
              [-O overhead] [-S sortby] [-U columns] [-a column] [-o file]
              [-s strsize] [-X format] [-P path]... [-p pid]...
              [--seccomp-bpf] { -p pid | [-DDD] [-E var[=val]]...
              [-u username] command [args] }
       strace -c [-dfwzZ] [-I n] [-b execve] [-e expr]... [-0 overhead]
              [-S sortby] [-U columns] [-P path]... [-p pid]...
              [--seccomp-bpf] { -p <u>pid</u> | [-DDD] [-E <u>var</u>[=<u>val</u>]]...
              [-u username] command [args] }
DESCRIPTION
       In the simplest case strace runs the specified command until it ex-
       its. It intercepts and records the system calls which are called by
       a process and the signals which are received by a process. The name
       of each system call, its arguments and its return value are printed
       on standard error or to the file specified with the -o option.
       strace is a useful diagnostic, instructional, and debugging tool.
       System administrators, diagnosticians and trouble-shooters will find
Manual page strace(1) line 1 (press h for help or q to quit)
```

\$ ls

_a.c a.out gdb-internals.txt \$ |



\$ man strace

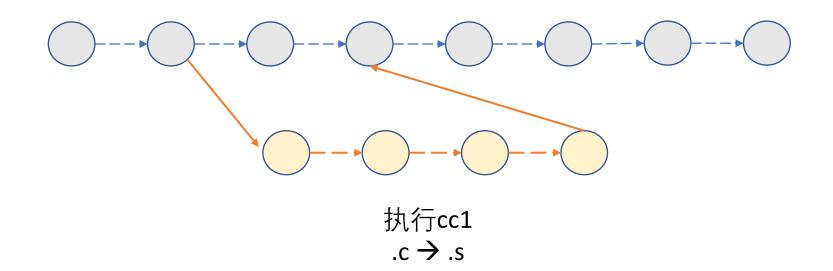
```
why@why-VirtualBox: ~/Documents/ICS2021/teach/Course17
                                                  why@why-VirtualBox: ~/Documents/ICS2021/teach/Course15
249 newfstatat(AT FDCWD, "/usr/lib/qcc/x86 64-linux-qnu/10/lto-wrapper", {st mode=S
    250 access("/usr/lib/gcc/x86 64-linux-gnu/10/lto-wrapper", X_0K) = 0
251 access("/tmp", R OK|W OK|X OK)
252 openat(AT FDCWD, "/tmp/ccFnrVKs.s", O RDWR|O CREAT|O EXCL, 0600) = 4
253 close(4)
254 newfstatat(AT_FDCWD, "/usr/lib/gcc/x86 64-linux-gnu/10/cc1", {st mode=S_IFREG|07
    55, st size=25697696, ...}, 0) = 0
255 access("/usr/lib/gcc/x86 64-linux-gnu/10/cc1", X OK) = 0
256 pipe2([4, 5], 0 CLOEXEC)
257 vfork( <unfinished ...>
258 [pid 17480] close(4)
259 [pid 17480] execve("/usr/lib/gcc/x86 64-linux-gnu/10/cc1", ["/usr/lib/gcc/x86 64
    -linux-gnu/10"..., "-quiet", "-imultiarch", "x86 64-linux-gnu", "a.c", "-quiet",
     "-dumpbase", "a.c", "-mtune=generic", "-march=x86-64", "-auxbase", "a", "-fasyn
    chronous-unwind-tables", "-fstack-protector-strong", "-Wformat", "-Wformat-secur
    ity", "-fstack-clash-protection", "-fcf-protection", "-o", "/tmp/ccFnrVKs.s"], 0
    x18ec000 /* 66 vars */ <unfinished ...>
260 [pid 17479] <... vfork resumed>)
                                             = 17480
261 [pid 17479] close(5)
                                             = 0
262 [pid 17479] read(4, "", 16)
263 [pid 17480] <... execve resumed>)
264 [pid 17479] close(4)
                                             = 0
265 [pid 17480] brk(NULL <unfinished ...>
266 [pid 17479] wait4(17480, <unfinished ...>
[No\ Name][+1]
                                               [strace] unix utf-8 Ln 266, Col 1/2923
```

364 [pid 17480] openat(AT_FDCWD, "/usr/lib/locale/locale-archive", O_RDONLY|O_CLOEXE

64 INFINITY) = 0

(C) = 4[No Name][+1]

.c到.s



262 [pid 17479] read(4, "", 16) = 6

263 [pid 17480] $< \dots$ execve resumed>) = 0

264 [pid 17479] close(4) = 0

265 [pid 17480] brk(NULL <unfinished ...>

[No Name][+1] [strace] unix utf-8 Ln 264, Col 1/2923

```
why@why-VirtualBox: ~/Documents/ICS2021/teach/Course17
                                                                        why@why-VirtualBox: ~/Documents/ICS2021/teach/Course15
1 \frac{e^{xecve}}{("/usr/lib/ccache/gcc", ["gcc", "a.c"], 0x7ffd5a558738 /* 60 vars */) = 0
```

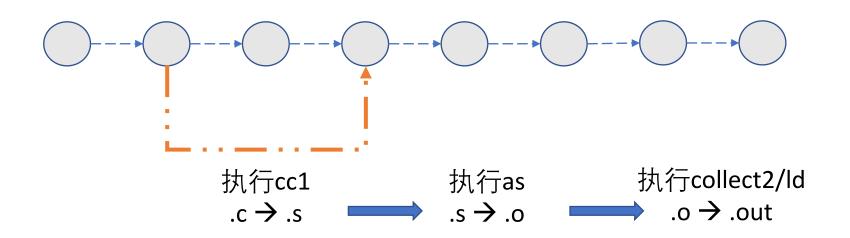
- 2 execve("/usr/bin/gcc", ["/usr/bin/gcc", "a.c"], 0x56492d0e5320 /* 61 vars */) = 0
- 3 [pid 17480] execve("/usr/lib/gcc/x86 64-linux-gnu/10/cc1", ["/usr/lib/gcc/x86 64linux-gnu/10"..., "-quiet", "-imultiarch", "x86 64-linux-gnu", "a.c", "-quiet", -dumpbase", "a.c", "-mtune=generic", "-march=x86-64", "-auxbase", "a", "-fasynchr
 - ond\$ echo \$PATH /usr/lib/ccache:/home/why/.local/bin:/usr/local/sbin:/usr/local/bin:/usr/sbin <code>jx18ec</code> [00]:/usr/bin:/sbin:/bin:/usr/games:/usr/local/games:/snap/bin:/snap/bin
- [pid 17480] <... execve resumed>)
- 5 [pid 17481] execve("/usr/lib/ccache/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or directo
- 6 [pid 17481] execve("/home/why/.local/bin/as", ["as", "--64", "-o", "/tmp/cc6160LT .o", "/tmp/ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or di rectory)
- 7 [pid 17481] execve("/usr/local/sbin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ccFnrVKs.s"], 0x18ec000 /* $6^{1}6$ vars */) = -1 ENOENT (No such file or directo
- 8 [pid 17481] execve("/usr/local/bin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", " /tmp/ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or director
- 9 [pid 17481] execve("/usr/sbin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or directory)
- 10 [pid 17481] execve("/usr/bin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/c cFnrVKs.s"], 0x18ec000 /* 66 vars */ <unfinished ...>

[strace] unix utf-8 Ln 1, Col 1/15 [No Name][+1]

- , ^
- 1 <mark>execve</mark>("/usr/lib/ccache/gcc", ["gcc", "a.c"], 0x7ffd5a558738 /* 60 vars */) = 0
- $2 \frac{\text{execve}}{\text{execve}}$ ("/usr/bin/gcc", "a.c"], 0x56492d0e5320 /* 61 vars */) = 0
- 3 [pid 17480] execve("/usr/lib/gcc/x86_64-linux-gnu/10/cc1", ["/usr/lib/gcc/x86_64-linux-gnu/10"..., "-quiet", "-imultiarch", "x86_64-linux-gnu", "a.c", "-quiet", "-dumpbase", "a.c", "-mtune=generic", "-march=x86-64", "-auxbase", "a", "-fasynchronous-unwind-tables", "-fstack-protector-strong", "-Wformat", "-Wformat-security", "-fstack-clash-protection", "-fcf-protection", "-o", "/tmp/ccFnrVKs.s"], 0x18ec 000 /* 66 vars */ <unfinished ...>
- $4 \text{ [pid } 17480 \text{]} < \dots \text{ execve resumed>)} = 0$
- 5 [pid 17481] $\frac{\text{execve}}{\text{ceche}}(\text{"/usr/lib/ccache/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ccFnrVKs.s"], <math>0x18ec000 /* 66 \text{ vars }*/) = -1 \text{ ENOENT (No such file or directory)}$
- 6 [pid 17481] execve("/home/why/.local/bin/as", ["as", "--64", "-o", "/tmp/cc6160LT
 .o", "/tmp/ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or di
 rectory)
- 7 [pid 17481] $\frac{\text{execve}}{\text{execve}}(\text{"/usr/local/sbin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ccFnrVKs.s"], <math>0x18ec000 \text{ /* } 6^{\frac{1}{10}} \text{ vars */}) = -1 \text{ ENOENT (No such file or directory)}$
- 8 [pid 17481] execve("/usr/local/bin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "
 /tmp/ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or director
 y)
- 9 [pid 17481] execve("/usr/sbin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ ccFnrVKs.s"], 0x18ec000 /* 66 vars */) = -1 ENOENT (No such file or directory)
- 10 [pid 17481] execve("/usr/bin/as", ["as", "--64", "-o", "/tmp/cc6160LT.o", "/tmp/ccFnrVKs.s"], 0x18ec000 /* 66 vars */ <unfinished ...>
- [No Name][+1] [strace] unix utf-8 Ln 1, Col 1/15

回到最初的问题

- 在IDE里,为什么按一个键,就能够编译运行?
 - 编译、链接
 - .c \rightarrow 预编译 \rightarrow .i \rightarrow 编译 \rightarrow .s \rightarrow 汇编 \rightarrow .o \rightarrow 链接 \rightarrow a.out
 - 加载执行
 - ./a.out
- 现在, 一位同学对这个过程提出了质疑
 - 如何说服这位同学?



答案: 用工具!

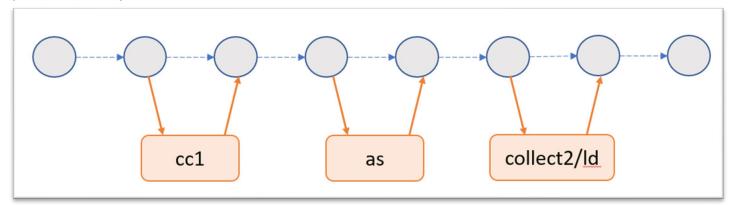
- 观察 trace: 查看 gcc 是否调用了其他命令
 - strace -f -qq gcc a.c 2>&1 | vim -
 - 可以使用 grep (shell) 或:g! (vim) 筛选感兴趣的系统调用
- 调试 gcc: 查看每一个阶段的中间结果
 - 在哪里打断点?
 - gdb-internals.txt

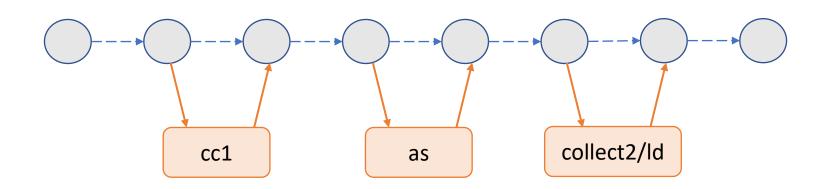


a.<u>c</u> a.out gdb-internals.txt

回到第一次C语言课

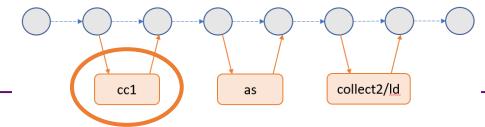
- 在IDE里,为什么按一个键,就能够编译运行?
 - 编译、链接
 - .c → 预编译 → .i → 编译 → .s → 汇编 → .o → 链接 → a.out
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 - 如何说服这位同学?



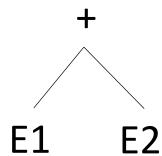


课堂上见过的轮子们

编译器 (.c → .s)



- 同 "表达式求值问题" (PA1)
 - 编译器 (OJ 题): 输入一个字符串, 输出计算它的汇编代码
 - 表达式 (a + b) * (c + d)
 - → 指令序列 (stack machine)
 - push \$a; push \$b; add; push \$c; push \$d; add; mul
- 现代编译器
 - 预编译 → 词法分析 → 语法树 → IR 中间代码 → 多趟优化 (内联、传播/ 折叠、删除冗余) → 指令生成/寄存器分配



汇编器 (.s → .o)

- cc1 as collect2/ld
- 除去预编译,汇编代码和指令几乎——对应
 - 根据指令集手册规定翻译
 - 生成符合 ELF 规范的二进制目标文件
 - printf("\x7f\x45\x4c\x46");...

```
      00000000:
      7f45
      4c46
      0201
      0100
      0000
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      00000
      00000
      00000</t
```

why@why-VirtualBox: ~/Documents/ICS2021/teach/Course17

_\$ ls
a.c a.out gdb-internals.txt

接器(.o)a.out)

- 按照 ld script 指示完成二进制文件构造、符号解析和重定位
 - gcc a.c -Wl,--verbose
 - "."代表当前位置;基本功能: 粘贴;定义符号

```
SECTIONS {
   = 0 \times 100000;
   .text : { *(entry) *(.text*) }
   etext = .; _etext = .;
   .rodata : { *(.rodata*) } .data : { *(.data*) }
   edata = .; _edata = .;
                                              .text
   .bss : { *(.bss*) }
   _start_start = ALIGN(4096);
                                              .data
                                                                .text
   . = _start_start + 0x8000;
   _stack_pointer = .;
                                              .text
                                                                .data
   end = .; _end = .;
                                              .data
   _{heap\_start} = ALIGN(4096);
   heap\_end = 0x8000000;
}
```

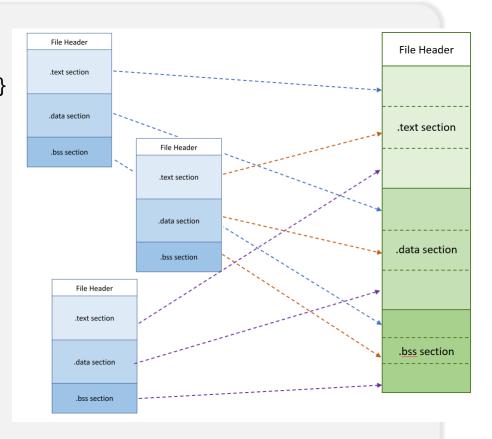
×

I

链接器 (.o → a.out)

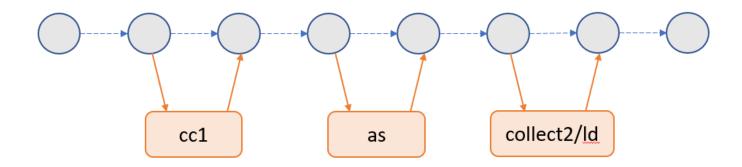
- 按照 ld script 指示完成二进制文件构造、符号解析和重定位
 - gcc a.c -Wl,--verbose
 - "."代表当前位置;基本功能: 粘贴;定义符号

```
SECTIONS {
   . = 0x100000;
   .text : { *(entry) *(.text*) }
   etext = .; _etext = .;
   .rodata : { *(.rodata*) }
   .data : { *(.data*) }
   edata = .; _edata = .;
   .bss : { *(.bss*) }
   _start_start = ALIGN(4096);
   . = _start_start + 0x8000;
   _stack_pointer = .;
   end = .; \_end = .;
   _{heap\_start} = ALIGN(4096);
   heap\_end = 0x8000000;
```



加载器

- PA里做过了
- 编译、链接、加载
 - 好像没那么难啊 (都是表达式求值和翻译)



单元测试框架

• YEMU 中的神奇代码

```
#define TESTCASES(_)
  \ _(1, 0b11100111, random_rm, ASSERT_EQ(newR[0], oldM[7]) )
  \ _(2, 0b00000100, random_rm, ASSERT_EQ(newR[1], oldR[0]) )
  \ _(3, 0b11100101, random_rm, ASSERT_EQ(newR[0], oldM[5]) )
  \ _(4, 0b00010001, random_rm, ASSERT_EQ(newR[0], oldR[0] + oldR[1]) )
  \ _(5, 0b11110111, random_rm, ASSERT_EQ(newM[7], oldR[0]) )
```

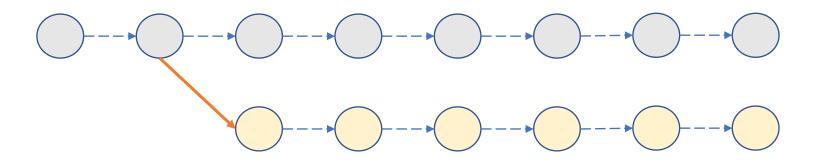
• 但是我相信大家都没有做好单元测试

Differential Testing

• 设置好状态; 各走一步

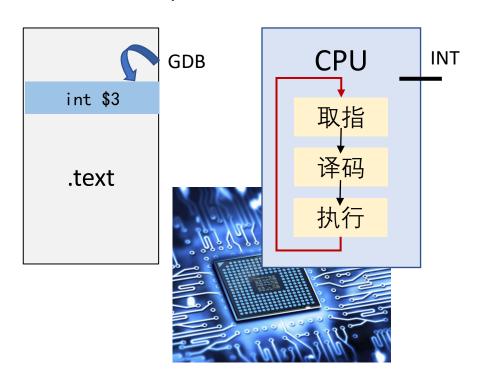
```
for i in range(int(sys.argv[1])):
    print('\n'.join(['si'] + [f'p ${r}' for r in ['eax', ...]]))
```

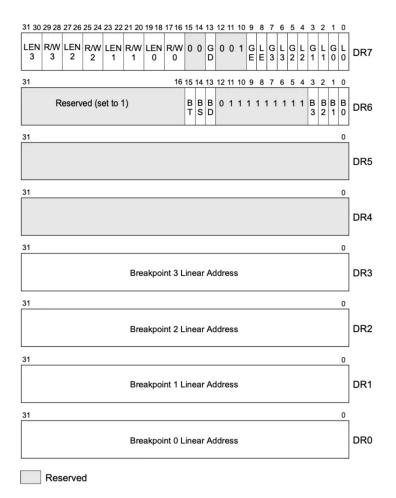
```
N=10000
    diff <(python3 cmdgen.py $N | ./x86-nemu-datui img) \
        <(python3 cmdgen.py $N | ./x86-nemu img)</pre>
```



调试器GDB

- 好奇它是怎么实现的?
- 考虑核心功能: 在任意 PC 的断点
- 其他功能都可以基于断点实现
 - 单步调试: 在下一条指令打断点
 - watch point: 单步调试 + 检查条件





why@why-VirtualBox: ~/Documents/ICS2021/teach/Course17

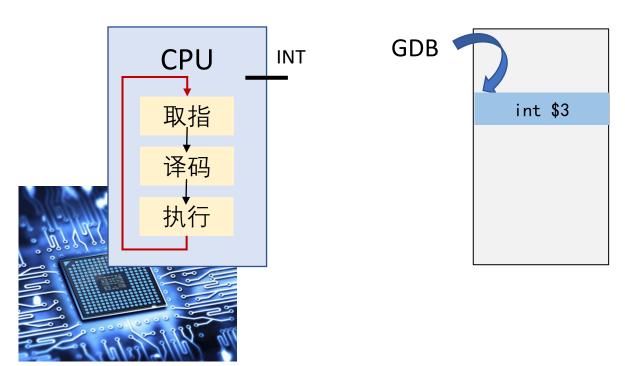
```
MPROTECT(2)
                           Linux Programmer's Manual
                                                                    MPROTECT(2)
NAME
      mprotect, pkey mprotect - set protection on a region of memory
SYNOPSIS
      #include <sys/mman.h>
       int mprotect(void *addr, size t len, int prot);
      #define GNU SOURCE
                                  /* See feature test macros(7) */
      #include <sys/mman.h>
      int pkey_mprotect(void *addr, size_t len, int prot, int pkey);
DESCRIPTION
      mprotect() changes the access protections for the calling process's memory
       pages containing any part of the address range in the
       [addr, addr+len-1]. addr must be aligned to a page boundary.
      If the calling process tries to access memory in a manner that violates
      the protections, then the kernel generates a SIGSEGV signal for the
      process.
      prot is a combination of the following access flags: PROT NONE or a bit-
      wise-or of the other values in the following list:
```

调试器GDB

- 好奇它是怎么实现的?
- 考虑核心功能
- 其他功能都可
 - 单步调试:

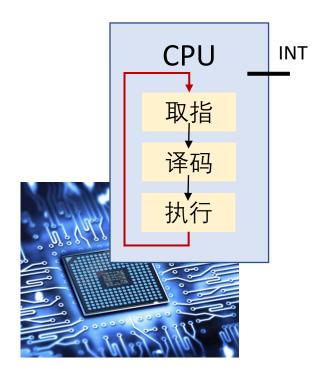
The **ptrace**() system call provides a means by which one process (the "tracer") may observe and control the execution of another process (the "tracee"), and examine and change the tracee's memory and registers. It is primarily used to implement breakpoint debugging and system call tracing.

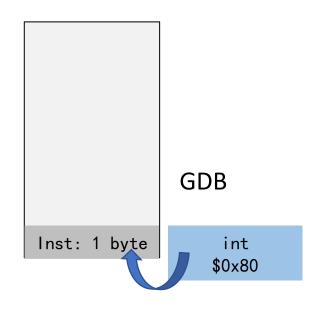
• watch point. ギシッツルマールロスコー



调试器GDB

- 好奇它是怎么实现的?
- 考虑核心功能: 在任意 PC 的断点
- 其他功能都可以基于断点实现
 - 单步调试: 在下一条指令打断点
 - watch point: 单步调试 + 检查条件



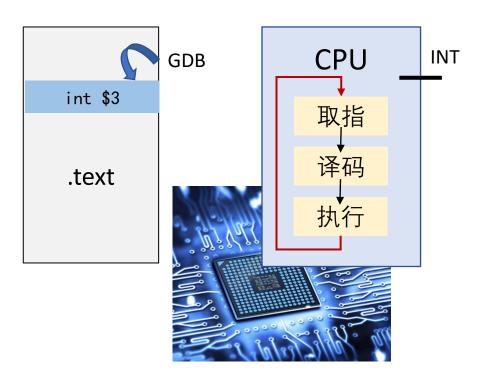


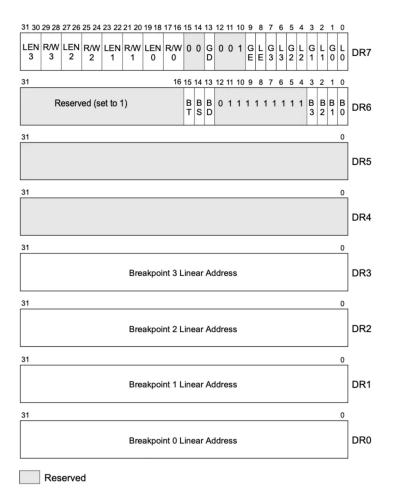
why@why-VirtualBox: ~/Documents/ICS2021/teach/Course17

\$ vim a.S

调试器GDB

- 好奇它是怎么实现的?
- 考虑核心功能: 在任意 PC 的断点
- 其他功能都可以基于断点实现
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 - watch point: 单步调试 + 检查条件





Trace/Profiler

- strace: 刚才已经展示过威力了
 - 如何实现?

```
$ strace ls
|execve("/usr/bin/ls", ["ls"], 0x7fff683feed0 /* 60 vars */) = 0
brk(NULL)
                                        = 0x55d1f69ba000
arch_prctl(0x3001 /* ARCH ??? */, 0x7fff7ba12c80) = -1 EINVAL (Invalid argument)
access("/etc/ld.so.preload", R OK)
                                       = -1 ENOENT (No such file or directory)
openat(AT FDCWD, "/etc/ld.so.cache", 0 RDONLY|0 CL0EXEC) = 3
|\text{newfstatat}(3, "", \{\text{st mode=S IFREG}|0644, st size=69006, ...}\}, AT EMPTY PATH) = 0
mmap(NULL, 69006, PROT READ, MAP PRIVATE, 3, 0) = 0 \times 7 \times 10^{-2}
close(3)
openat(AT FDCWD, "/lib/x86 64-linux-gnu/libselinux.so.1", 0 RDONLY|0 CLOEXEC) = 3
|\text{newfstatat}(3, "", \{\text{st mode=S IFREG}|0644, st size=166272, ...}, AT EMPTY PATH) = 0
mmap(NULL, 8192, PROT READ|PROT WRITE, MAP PRIVATE|MAP ANONYMOUS, -1, 0) = 0 \times 7 \text{ ff} 23 \text{ co}
90000
mmap(NULL, 177672, PROT READ, MAP PRIVATE | MAP DENYWRITE, 3, 0) = 0x7ff23c064000
mmap(0x7ff23c06a000, 106496, PROT READ|PROT EXEC, MAP PRIVATE|MAP FIXED|MAP DENYWRIT
E, 3, 0x6000) = 0x7ff23c06a000
mmap(0x7ff23c084000, 32768, PROT READ, MAP PRIVATE|MAP FIXED|MAP DENYWRITE, 3, 0x200
|00\rangle = 0x7ff23c084000
mmap(0x7ff23c08c000, 8192, PROT READ|PROT WRITE, MAP PRIVATE|MAP FIXED|MAP DENYWRITE
3, 0x27000) = 0x7ff23c08c000
mmap(0x7ff23c08e000, 5640, PROT READ|PROT WRITE, MAP PRIVATE|MAP FIXED|MAP ANONYMOUS
(-1, 0) = 0x7ff23c08e000
close(3)
openat(AT FDCWD, "/lib/x86 64-linux-gnu/libc.so.6", 0 RDONLY 0 CLOEXEC) = 3
read(3, "\177ELF\2\1\1\3\0\0\0\0\0\0\0\0\0\1\0\0\0\240\206\2\0\0\0\0\0"..., 832
 = 832
```

PTRACE(2)

Linux Programmer's Manual

PTRACE(2)

NAME

```
ptrace - process trace
```

SYNOPSIS

```
#include <sys/ptrace.h>
```

DESCRIPTION

The **ptrace**() system call provides a means by which one process (the "tracer") may observe and control the execution of another process (the "tracee"), and examine and change the tracee's memory and registers. It is primarily used to implement breakpoint debugging and system call tracing.

A tracee first needs to be attached to the tracer. Attachment and subsequent commands are per thread: in a multithreaded process, every thread can be individually attached to a (potentially different) tracer, or left not attached and thus not debugged. Therefore, "tracee" always means "(one) thread", never "a (possibly multithreaded) process". Ptrace commands are always sent to a specific tracee using a call of the form

```
ptrace(PTRACE foo, pid, ...)
```

Manual page ptrace(2) line 1 (press h for help or q to quit)

静态分析工具 Lint

- 大家见到的第一个 lint: gcc -Wall -Werror
 - cppcheck
 - Cert C Coding Standard
 - 自带 <u>checker</u>

Introduction

由 Robert Seacord创建, 最终由 Robert Schiela修改于 一月 03, 2017

- Scope
- Audience
- How this Coding Standard is Organized
- History
- ISO/IEC TS 17961 C Secure Coding Rules
- Tool Selection and Validation
- Taint Analysis
- Rules versus Recommendations
- Conformance Testing
- Development Process
- Usage
- System Qualities
- Automatically Generated Code
- Government Regulations
- Acknowledgments

This standard provides rules for secure coding in the C programming language. The goal of these rules and recommendations is to develop safe, reliable, and secure systems, for example by eliminating undefined behaviors that can lead to undefined program behaviors and exploitable vulnerabilities. Conformance to the coding rules defined in this standard are necessary (but not sufficient) to ensure the safety, reliability, and security of software systems developed in the C programming language. It is also necessary, for example, to have a safe and secure design. Safety-critical systems typically have stricter requirements than are imposed by this coding standard, for example requiring that all memory be statically allocated. However, the application of this coding standard will result in high-quality systems that are reliable, robust, and resistant to attack.

动态分析工具 Sanitizer

- 本质:运行时额外增加的 assert()
 - 回顾调试理论

- 一些非常实用的 assertions
 - assert(IS_GUESTPTR(ptr));
 - (PMEM_MAP_START <= (x) && (x) < PMEM_MAP_END)
 - assert(IS SMALLINT(x));
 - $(0 \le (x) \&\& (x) \le 100)$
 - assert(IS_BOOL(ptr));
 - ((x) == 0 | | (x) == 1)

总结

我们学到了什么?

- "程序是个状态机。"
- "我们可以观测状态机的设计、实现和执行。"
- 能实现几乎任何工具(的玩具版)
 - 并且能在需要的使用善用它们
 - 编译器 (gcc)
 - 汇编器 (as)
 - 链接器 (ld)
 - 调试器 (gdb)
 - 追踪器 (strace/ltrace)
 - profiler (perf)

Cheers!

(你们完成了了不起的事情!)